A Proof of Lemma 1

Proof (Lemma 1).

$$\hat{L}(q; \mathcal{S}) - \hat{L}(q'; \mathcal{S}) = \hat{E}_{\mathcal{S}}[\log q'(x) - \log q(x)]$$

$$= \hat{E}_{\mathcal{S}} \left[-af(x) - \log Z_{q'}\right]$$

$$= a(\mathbb{E}_q \left[f(x)\right] - \hat{\mathbb{E}}_{\mathcal{S}} \left[f(x)\right]) - a\mathbb{E}_q \left[f(x)\right] - \log Z_{q'}$$

$$= a\hat{\alpha}(f) - a\mathbb{E}_q \left[f(x)\right] - \log Z_{q'}$$
(4)

Since $\hat{\alpha}(f) \ge a$ by assumption, it remains to show $a \mathbb{E}_q [f(x)] + \log Z_{q'} \le a^2/2$. Using the bound $\log r \le r - 1$ for any r > 0, we get that,

$$\begin{aligned} a \mathbf{E}_q \left[f(x) \right] + \log Z_{q'} &\leq a \mathbf{E}_q \left[f(x) \right] + Z_{q'} - 1 \\ &= a \mathbf{E}_q \left[f(x) \right] + \mathbf{E}_q [e^{-af(x)}] - 1 \\ &= \mathbf{E}_q \left[af(x) + e^{-af(x)} - 1 \right] \\ &\leq \mathbf{E}_q [(af(x))^2/2], \end{aligned}$$

where we have used the fact that $Z_{q'} = E_q[e^{-af(x)}]$ and, to get to the last line we use $e^{-r} + r - 1 \le r^2/2$ for $r \ge 0$ by Taylor expansion. Since $f(x) \in [0, 1]$, the last quantity is at most $a^2/2$, which together with (4), gives $\hat{L}(q; S) - \hat{L}(q'; S) \ge a^2/2$.

B Proof of Lemma 2

Proof (Lemma 2). Let $g(x) = \log q(x) - \log q'(x)$. By Jensen's inequality,

$$\begin{aligned} \mathbf{E}_q[g(x)] - \hat{\mathbf{E}}_{\mathcal{S}}[g(x)] &= -\mathbf{E}_q \left[\log \frac{q'(x)}{q(x)} \right] - \hat{\mathbf{E}}_{\mathcal{S}}[g(x)] \\ &\geq -\log \mathbf{E}_q \left[\frac{q'(x)}{q(x)} \right] - \hat{\mathbf{E}}_{\mathcal{S}}[g(x)] \\ &= -\log(1) - \hat{\mathbf{E}}_{\mathcal{S}}[g(x)] \\ &= \mathbf{E}_{\mathcal{S}}[\log q'(x)] - \hat{\mathbf{E}}_{\mathcal{S}}[\log q(x)] \\ &= \hat{L}(q;S) - \hat{L}(q';S) \end{aligned}$$

Since $f(x) = \frac{1}{2 \log C}(g(x) + \log C)$, the training advantage of f is that of g scaled by a factor of $\frac{1}{2 \log C}$. Finally, it is straightforward to verify that $f(x) \in [0, 1]$ by our assumptions on the ratio between q and q'.

C Proof of Lemma 3

Proof. We proceed analogously as in the proof of Lemma 1. We first note that

$$\hat{L}(q; S) = -\hat{E}_{S} \left[\log \prod_{i=1}^{N} q(x_{i} \mid x_{1}, \dots, x_{i-1}) \right] = -\sum_{i=1}^{N} \hat{E}_{S} \log q(x_{i} \mid x_{1}, \dots, x_{i-1}),$$

and

$$\hat{L}(q'; S) = -\hat{E}_{S} \left[\log \prod_{i=1}^{N} q'(x_{i} \mid x_{1}, \dots, x_{i-1}) \right]$$
$$= \sum_{i=1}^{N} -\hat{E}_{S} \log q(x_{i} \mid x_{1}, \dots, x_{i-1}) + bg(x_{1}, \dots, x_{i}) + \log Z_{q'}(x_{1}, \dots, x_{i-1})$$

Let us use the short-hand notation $x_{1:i} \triangleq (x_1, \ldots, x_i)$. Subtracting the two equalities above we obtain

$$\hat{L}(q; \mathcal{S}) - \hat{L}(q'; \mathcal{S}) = \sum_{i=1}^{N} \hat{E}_{\mathcal{S}} \left[-bg(x_{1:i}) - \log Z_{q'}(x_{1:i-1}) \right],$$

which, after adding and subtracting $\hat{E}_{\mathcal{S}} E_{w \sim q(w|x_{1:i-1})} g(x_{1:i-1}, w)$ and rearranging terms, yields

$$\hat{L}(q; \mathcal{S}) - \hat{L}(q'; \mathcal{S}) = b \left[\sum_{i=1}^{N} \hat{E}_{\mathcal{S}} \left(E_{w \sim q(w|x_{1:i})} g(x_{1:i-1}, w) - g(x_{1:i}) \right) \right]$$
(5)

$$-\sum_{i=1}^{N} \hat{\mathbf{E}}_{\mathcal{S}} \left[\log Z_{q'}(x_{1:i-1}) - b \mathbf{E}_{w \sim q(w|x_{1:i-1})} g(x_{1:i-1}, w) \right]$$
(6)

$$=bN\hat{\beta}(g) - \sum_{i=1}^{N} \hat{\mathbf{E}}_{\mathcal{S}} \left[b\mathbf{E}_{w}g(x_{1:i-1}, w) + \log Z_{q'}(x_{1:i-1})\right]$$
(7)

By assumption we have $Nb\hat{\beta}(g) \ge Nb^2$, so it it remains to show that the second term is upper bounded by $Nb^2/2$. Using, as before, the bound $\log r \le r-1$ for every $r = Z_{q'}(x_{1:i-1}) \ge 0$, we get that, for every $i = 1, \ldots, N$:

$$\begin{split} \hat{\mathbf{E}}_{\mathcal{S}} \left[b \mathbf{E}_{w} g(x_{1:i-1}, w) + \log Z_{q'}(x_{1:i-1}) \right] &\leq \hat{\mathbf{E}}_{\mathcal{S}} \left[b \mathbf{E}_{w} g(x_{1:i-1}, w) + Z_{q'}(x_{1:i-1}) - 1 \right] \\ &= \hat{\mathbf{E}}_{\mathcal{S}} \left[b \mathbf{E}_{w} g(x_{1:i-1}, w) + \mathbf{E}_{w} e^{-bg(x_{1:i-1}, w)} - 1 \right] \\ &= \hat{\mathbf{E}}_{\mathcal{S}} \mathbf{E}_{w} \left[bg(x_{1:i-1}, w) + e^{-bg(x_{1:i-1}, w)} - 1 \right] \\ &\leq \hat{\mathbf{E}}_{\mathcal{S}} \mathbf{E}_{w} \left(bg(x_{1:i-1}, w)/2 \right)^{2} \leq \frac{b^{2}}{2} \end{split}$$

where the last inequality follows again from the fact that $g(x) \in [0, 1]$ for any x. Therefore, the sum over these N terms is upper bounded by $N\frac{b^2}{2}$, which combined with (7), yields the desired result. \Box

D Proof of Theorem 1

Proof. The fact that Algorithm 1 terminates with a distribution q which is ϵ -indistinguishable by O_d is immediate from the stopping criterion.

Now, for the runtime analysis, note that —by construction— the iterates $g_t, t \in \{0, \ldots, T-1\}$ have training advantage $\hat{\beta}(g_t, S, q_t) \geq \epsilon$. Thus, by Lemma 3, the algorithm makes at least $\frac{N\epsilon^2}{2}$ improvement in each iteration. Therefore, the total number of iterations T is at most $\frac{2L_0}{N\epsilon^2}$, where $L_0 := \hat{L}(q_0; S)$ is the log-loss of the initial model. Each iteration of Algorithm 1 requires calling O_d oracle once, evaluating $\hat{\beta}(\cdot)$ at an $\mathcal{O}(NmnT_g)$ complexity, and updating each of the n next-token probabilities of q for each sequence length $1, \ldots, N$. Each of these updates involves evaluating g plus an $\mathcal{O}(n)$ partition normalization. Putting these together, we conclude that each iteration has $O(T_d + NnT_g(m + n))$ complexity.

Combining the the two arguments above, we conclude that Algorithm 1 has a total runtime of $O(\frac{1}{\epsilon^2}L_0(\frac{T_d}{N} + nT_g(m+n)))$.

E Empirical validation

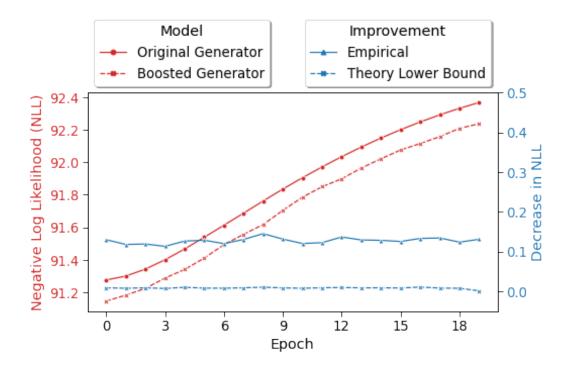


Figure 1: **Empirical Validation of Lemma 3**. For a simple pre-trained SeqGAN model (generator + discriminator), we show that the boosting scheme proposed in that Lemma results in reduced log-loss (NLL) throughout training. Furthermore, the empirical difference between original and boosted models is indeed lower-bounded by the gap predicted by the Lemma.