

A Proof of query complexities

A.1 Properties of adaptive sequential testing in Procedure 2

Lemma 1. Suppose $\{X_i\}_{i=1}^\infty$ is a sequence of i.i.d. random variables such that $\mathbb{E}X_i \leq 0$, $|X_i| \leq 1$. Let $\delta > 0$. Then with probability at least $1 - \delta$, for all $n \in \mathbb{N}$ simultaneously $\text{CheckSignificant}(\{X_i\}_{i=1}^n, \delta)$ in Procedure 2 returns false.

Proof. This is immediate by applying Proposition 1 to $X_i - \mathbb{E}X_i$. \square

Lemma 2. Suppose $\{X_i\}_{i=1}^\infty$ is a sequence of i.i.d. random variables such that $\mathbb{E}X_i > \epsilon > 0$, $|X_i| \leq 1$. Let $\delta \in [0, \frac{1}{3}]$, $N \geq \frac{\xi}{\epsilon^2} \ln \frac{1}{\delta} [\ln \ln]_+ \frac{1}{\epsilon}$ (ξ is an absolute constant specified in the proof). Then with probability at least $1 - \delta$, $\text{CheckSignificant}(\{X_i\}_{i=1}^N, \delta)$ in Procedure 2 returns true.

Proof. Let $S_N = \sum_{i=1}^N X_i$. $\text{CheckSignificant}(\{X_i\}_{i=1}^N, \delta)$ returns false if and only if $S_N \leq D_0 \left(1 + \ln \frac{1}{\delta} + \sqrt{N([\ln \ln]_+ N + \ln \frac{1}{\delta})}\right)$.

$$\begin{aligned} & \Pr \left(S_N \leq D_0 \left(1 + \ln \frac{1}{\delta} + \sqrt{N([\ln \ln]_+ N + \ln \frac{1}{\delta})}\right) \right) \\ & \leq \Pr \left(S_N \leq D_0 \left(1 + \ln \frac{1}{\delta} + \sqrt{N[\ln \ln]_+ N} + \sqrt{N \ln \frac{1}{\delta}}\right) \right) \\ & \leq \Pr \left(S_N - N\mathbb{E}X_i \leq D_0 \left(1 + \ln \frac{1}{\delta} + \sqrt{N[\ln \ln]_+ N} + \sqrt{N \ln \frac{1}{\delta}}\right) - N\epsilon \right) \end{aligned}$$

Suppose $N = \frac{c\xi}{\epsilon^2} \ln \frac{1}{\delta} [\ln \ln]_+ \frac{1}{\epsilon}$ for constant $c \geq 1$ and ξ . ξ is set to be sufficiently large, such that (1) $\xi \geq 4D_0^2$; (2) $\frac{2D_0}{\sqrt{\xi}} + D_0 \left(3 + \sqrt{[\ln \ln]_+ \xi}\right) + D_0 - \sqrt{\xi}/2 \leq -\sqrt{\frac{1}{2}}$; (3) $f(x) = D_0 \sqrt{[\ln \ln]_+ x} - \sqrt{x}/2$ is decreasing when $x > \xi$. Here (2) is satisfiable since $\frac{D_0}{\sqrt{\xi}} + D_0 \sqrt{[\ln \ln]_+ \xi} - \sqrt{\xi}/2 \rightarrow -\infty$ as $\xi \rightarrow \infty$, (3) is satisfiable since $f'(x) \rightarrow -\infty$ as $x \rightarrow \infty$. (2) and (3) together implies $\frac{2D_0}{\sqrt{\xi}} + D_0 \left(3 + \sqrt{[\ln \ln]_+ c\xi}\right) + D_0 - \sqrt{c\xi}/2 \leq -\sqrt{\frac{1}{2}}$.

$$\begin{aligned} & \frac{1}{\sqrt{N}} \left(D_0 \left(1 + \ln \frac{1}{\delta} + \sqrt{N[\ln \ln]_+ N} + \sqrt{N \ln \frac{1}{\delta}}\right) - N\epsilon \right) \\ & = \sqrt{\ln \frac{1}{\delta}} \left(\frac{D_0 \epsilon (1 + \ln \frac{1}{\delta})}{\sqrt{c\xi [\ln \ln]_+ \frac{1}{\epsilon} \ln \frac{1}{\delta}}} + D_0 \sqrt{\frac{[\ln \ln]_+ \left(\frac{c\xi}{\epsilon^2} \ln \frac{1}{\delta} [\ln \ln]_+ \frac{1}{\epsilon} \right)}{\ln \frac{1}{\delta}}} + D_0 - \sqrt{c\xi [\ln \ln]_+ \frac{1}{\epsilon}} \right) \end{aligned}$$

Since $[\ln \ln]_+ \frac{1}{\epsilon}, c, \ln \frac{1}{\delta} \geq 1$ and $\epsilon < 1$, we have $\frac{D_0 \epsilon (1 + \ln \frac{1}{\delta})}{\sqrt{c\xi [\ln \ln]_+ \frac{1}{\epsilon} \ln \frac{1}{\delta}}} \leq \frac{2D_0}{\sqrt{\xi}}$.

Since $[\ln \ln]_+ x \geq 1$ if $x \geq 1$, we have $[\ln \ln]_+ \frac{1}{\epsilon} \leq \frac{1}{\epsilon}$, and thus

$$\begin{aligned}
\sqrt{[\ln \ln]_+ \left(\frac{c\xi}{\epsilon^2} \ln \frac{1}{\delta} [\ln \ln]_+ \frac{1}{\epsilon} \right)} &= \sqrt{\ln \left[\max \left\{ e, 2 \ln \frac{1}{\epsilon} + \ln c\xi + \ln \ln \frac{1}{\delta} + \ln [\ln \ln]_+ \frac{1}{\epsilon} \right\} \right]} \\
&\leq \sqrt{\ln \left[\max \left\{ e, 3 \ln \frac{1}{\epsilon} + \ln c\xi + [\ln \ln]_+ \frac{1}{\delta} \right\} \right]} \\
&\stackrel{(a)}{\leq} \sqrt{\ln \left[\max \left\{ e, 9 \ln \frac{1}{\epsilon} \ln c\xi [\ln \ln]_+ \frac{1}{\delta} \right\} \right]} \\
&\leq \sqrt{3 + [\ln \ln]_+ \frac{1}{\epsilon} + [\ln \ln]_+ c\xi + \ln [\ln \ln]_+ \frac{1}{\delta}} \\
&\stackrel{(b)}{\leq} \sqrt{3} + \sqrt{[\ln \ln]_+ c\xi} + \sqrt{[\ln \ln]_+ \frac{1}{\epsilon}} + \sqrt{\ln [\ln \ln]_+ \frac{1}{\delta}}
\end{aligned}$$

where (a) follows by $a + b + c \leq 3abc$ if $a, b, c \geq 1$, and (b) follows by $\sqrt{\sum_i x_i} \leq \sum_i \sqrt{x_i}$ if $x_i \geq 0$.

Thus, we have

$$\begin{aligned}
&\frac{1}{\sqrt{N}} \left(D_0 \left(1 + \ln \frac{1}{\delta} + \sqrt{N[\ln \ln]_+ N} + \sqrt{N \ln \frac{1}{\delta}} \right) - N\epsilon \right) \\
&\leq \sqrt{\ln \frac{1}{\delta}} \left(\frac{2D_0}{\sqrt{\xi}} + D_0 \frac{\sqrt{3} + \sqrt{[\ln \ln]_+ c\xi} + \sqrt{[\ln \ln]_+ \frac{1}{\epsilon}} + \sqrt{\ln [\ln \ln]_+ \frac{1}{\delta}}}{\sqrt{\ln \frac{1}{\delta}}} + D_0 - \sqrt{c\xi [\ln \ln]_+ \frac{1}{\epsilon}} \right) \\
&\stackrel{(c)}{\leq} \sqrt{\ln \frac{1}{\delta}} \left(\frac{2D_0}{\sqrt{\xi}} + D_0 (3 + \sqrt{[\ln \ln]_+ c\xi}) + D_0 - \sqrt{c\xi/2} \right) \\
&\stackrel{(d)}{\leq} -\sqrt{\ln \frac{1}{\delta}}/2
\end{aligned}$$

(c) follows by $\sqrt{\ln \frac{1}{\delta}} \geq \max \left\{ 1, \sqrt{\ln [\ln \ln]_+ \frac{1}{\delta}} \right\}$, $D_0 \geq 1$, and $\sqrt{[\ln \ln]_+ \frac{1}{\epsilon}} \left(\frac{D_0}{\sqrt{\ln \frac{1}{\delta}}} - \sqrt{c\xi} \right) \leq D_0 - \sqrt{c\xi} \leq -\sqrt{c\xi/2}$ if $c\xi \geq 4D_0^2$. (d) follows by our choice of ξ .

Therefore,

$$\begin{aligned}
&\Pr \left(S_N - N\mathbb{E}X_i \leq D_0 \left(1 + \ln \frac{1}{\delta} + \sqrt{N[\ln \ln]_+ N} + \sqrt{N \ln \frac{1}{\delta}} \right) - N\epsilon \right) \\
&\leq \Pr \left(S_N - N\mathbb{E}X_i \leq -\sqrt{N \ln \frac{1}{\delta}}/2 \right)
\end{aligned}$$

which is at most δ by Hoeffding Bound. \square

Lemma 3. Suppose $\{X_i\}_{i=1}^\infty$ is a sequence of i.i.d. random variables such that $\mathbb{E}X_i \leq 0$, $|X_i| \leq 1$. Let $\delta > 0$. Then with probability at least $1 - \delta$, for all n simultaneously CheckSignificant-Var($\{X_i\}_{i=1}^n, \delta$) in Procedure 2 returns false.

Proof. Define $Y_i = X_i - \mathbb{E}X_i$. It is easy to check $\frac{n}{n-1} \left(\sum_{i=1}^n Y_i^2 - \frac{1}{n} (\sum_{i=1}^n Y_i)^2 \right) = \frac{n}{n-1} \left(\sum_{i=1}^n X_i^2 - \frac{1}{n} (\sum_{i=1}^n X_i)^2 \right)$. The result is immediate from Proposition 2. \square

Lemma 4. Suppose $\{X_i\}_{i=1}^\infty$ is a sequence of i.i.d. random variables such that $\mathbb{E}X_i > \tau\epsilon$, $|X_i| \leq 1$, $\text{Var}(X_i) \leq 2\epsilon$ where $0 < \epsilon \leq 1$, $\tau > 0$. Let $\delta < 1$, $N = \frac{\xi}{\tau\epsilon} \ln \frac{2}{\delta}$ (ξ is a constant specified in the proof). Then with probability at least $1 - \delta$, $\text{CheckSignificant-Var}(\{X_i\}_{i=1}^N, \delta)$ in Procedure 2 returns true.

Proof. Let $Y_i = X_i - \mathbb{E}X_i$, η be the constant η in Lemma 14. Set $\xi = \max(\eta, \frac{16}{\tau} + \frac{8}{3})$.

$\text{CheckSignificant-Var}(\{X_i\}_{i=1}^N, \delta)$ returns false if and only if $\sum_{i=1}^N X_i \leq q(N, \text{Var}, \delta)$.

By applying Lemma 14 to X_i , $\frac{q(N, \text{Var}, \delta)}{N} - \mathbb{E}X_i \leq -\tau\epsilon/2$ with probability at least $1 - \delta/2$.

Applying Bernstein's inequality to Y_i , we have

$$\begin{aligned} \Pr\left(\frac{1}{N} \sum_{i=1}^N Y_i \leq -\tau\epsilon/2\right) &\leq \exp\left(-\frac{N(-\tau\epsilon)^2/4}{4\epsilon + 2\tau\epsilon/3}\right) \\ &= \exp\left(-\frac{\xi \ln \frac{2}{\delta}}{16/\tau + 8/3}\right) \\ &\leq \delta/2 \end{aligned}$$

Thus, by a union bound,

$$\begin{aligned} &\Pr\left(\sum_{i=1}^N X_i \leq q(N, \text{Var}, \delta)\right) \\ &\leq \Pr\left(\frac{q(N, \text{Var}, \delta)}{N} - \mathbb{E}X_i \geq -\tau\epsilon/2\right) \\ &\quad + \Pr\left(\frac{q(N, \text{Var}, \delta)}{N} - \mathbb{E}X_i \leq -\tau\epsilon/2 \text{ and } \frac{1}{N} \sum_{i=1}^N X_i \leq \frac{q(N, \text{Var}, \delta)}{N}\right) \\ &\leq \delta/2 + \Pr\left(\frac{q(N, \text{Var}, \delta)}{N} - \mathbb{E}X_i \leq -\tau\epsilon/2 \text{ and } \frac{1}{N} \sum_{i=1}^N Y_i \leq \frac{q(n, \text{Var}, \delta)}{N} - \mathbb{E}X_i\right) \\ &\leq \delta/2 + \Pr\left(\frac{1}{N} \sum_{i=1}^N Y_i \leq -\tau\epsilon/2\right) \\ &\leq \delta \end{aligned}$$

□

A.2 The one-dimensional case

Proof of Theorem 1. Since $\hat{\theta} = (L_{\log \frac{1}{2\epsilon}} + R_{\log \frac{1}{2\epsilon}})/2$ and $R_{\log \frac{1}{2\epsilon}} - L_{\log \frac{1}{2\epsilon}} = 2\epsilon$, $|\hat{\theta} - \theta^*| > \epsilon$ is equivalent to $\theta^* \notin [L_{\log \frac{1}{2\epsilon}}, R_{\log \frac{1}{2\epsilon}}]$. We have

$$\begin{aligned} \Pr(|\hat{\theta} - \theta^*| > \epsilon) &= \Pr(\theta^* \notin [L_{\log \frac{1}{2\epsilon}}, R_{\log \frac{1}{2\epsilon}}]) \\ &= \Pr(\exists k : \theta^* \in [L_k, R_k] \text{ and } \theta^* \notin [L_{k+1}, R_{k+1}]) \\ &\leq \sum_{k=0}^{\log \frac{1}{2\epsilon} - 1} \Pr(\theta^* \in [L_k, R_k] \text{ and } \theta^* \notin [L_{k+1}, R_{k+1}]) \end{aligned}$$

For any $k = 0, \dots, \log \frac{1}{2\epsilon} - 1$, define $\mathbb{Q}_k = \left\{ (p, q) : p, q \in \mathbb{Q} \cap [0, 1] \text{ and } q - p = \left(\frac{3}{4}\right)^k \right\}$ where \mathbb{Q} is the set of rational numbers. Note that $L_k, R_k \in \mathbb{Q}_k$, and \mathbb{Q} is countable. So we have

$$\begin{aligned}
& \Pr(\theta^* \in [L_k, R_k] \text{ and } \theta^* \notin [L_{k+1}, R_{k+1}]) \\
&= \sum_{(p,q) \in \mathbb{Q}_k: p \leq \theta^* \leq q} \Pr(L_k = p, R_k = q \text{ and } \theta^* \notin [L_{k+1}, R_{k+1}]) \\
&= \sum_{(p,q) \in \mathbb{Q}_k: p \leq \theta^* \leq q} \Pr(\theta^* \notin [L_{k+1}, R_{k+1}] | L_k = p, R_k = q) \Pr(L_k = p, R_k = q)
\end{aligned}$$

Define event $E_{k,p,q}$ to be the event $L_k = p, R_k = q$. To show $\Pr(|\hat{\theta} - \theta^*| > \epsilon) \leq \frac{\delta}{2}$, it suffices to show $\Pr(\theta^* \notin [L_{k+1}, R_{k+1}] | E_{k,p,q}) \leq \frac{\delta}{2 \log \frac{1}{2\epsilon}}$ for any $k = 0, \dots, \log \frac{1}{2\epsilon} - 1$, $(p, q) \in \mathbb{Q}_k$ and $p \leq \theta^* \leq q$.

Conditioning on event $E_{k,p,q}$, event $\theta^* \notin [L_{k+1}, R_{k+1}]$ happens only if some calls of CheckSignificant and CheckSignificant-Var between Line 16 and 27 of Algorithm 1 return true incorrectly. In other words, at least one of following events happens for some n :

- $O_{k,p,q}^{(1)}: \theta^* \in [L_k, U_k]$ and CheckSignificant-Var($\{A_i^{(u)} - A_i^{(m)}\}_{i=1}^n, \frac{\delta}{4 \log \frac{1}{2\epsilon}}$) returns true;
- $O_{k,p,q}^{(2)}: \theta^* \in [V_k, R_k]$ and CheckSignificant-Var($\{A_i^{(v)} - A_i^{(m)}\}_{i=1}^n, \frac{\delta}{4 \log \frac{1}{2\epsilon}}$) returns true;
- $O_{k,p,q}^{(3)}: \theta^* \in [L_k, U_k]$ and CheckSignificant($\{-B_i^{(u)}\}_{i=1}^n, \frac{\delta}{4 \log \frac{1}{2\epsilon}}$) returns true;
- $O_{k,p,q}^{(4)}: \theta^* \in [V_k, R_k]$ and CheckSignificant($\{B_i^{(v)}\}_{i=1}^n, \frac{\delta}{4 \log \frac{1}{2\epsilon}}$) returns true;

Note that since $[U_k, V_k] \subset [L_{k+1}, R_{k+1}]$ for any k by our construction, if $\theta^* \in [U_k, V_k]$ then $\theta^* \in [L_{k+1}, R_{k+1}]$. Besides, event $\theta^* \in [L_k, U_k]$ and event $\theta^* \in [V_k, R_k]$ are mutually exclusive.

Conditioning on event $E_{k,p,q}$, suppose for now $\theta^* \in [L_k, U_k]$.

$$\begin{aligned}
& \Pr(O_{k,p,q}^{(1)} | E_{k,p,q}) \\
&= \Pr\left(\exists n : \text{CheckSignificant-Var}\left(\{D_i^{(u,m)}\}_{i=1}^n, \frac{\delta}{4 \log \frac{1}{2\epsilon}}\right) \text{ returns true} \mid \theta^* \in [L_k, U_k], E_{k,p,q}\right)
\end{aligned}$$

On event $\theta^* \in [L_k, U_k]$ and $E_{k,p,q}$, the sequences $\{A_i^{(u)}\}$ and $\{A_i^{(m)}\}$ are i.i.d., and $\mathbb{E}[A_i^{(u)} - A_i^{(m)} \mid \theta^* \in [L_k, U_k], E_{k,p,q}] \leq 0$. By Lemma 3, the probability above is at most $\frac{\delta}{4 \log \frac{1}{2\epsilon}}$.

Likewise,

$$\begin{aligned}
& \Pr(O_{k,p,q}^{(3)} | E_{k,p,q}) \\
&= \Pr\left(\exists n : \text{CheckSignificant}\left(\{-B_i^{(u)}\}_{i=1}^n, \frac{\delta}{4 \log \frac{1}{2\epsilon}}\right) \text{ returns true} \mid \theta^* \in [L_k, U_k], E_{k,p,q}\right)
\end{aligned}$$

On event $\theta^* \in [L_k, U_k]$ and $E_{k,p,q}$, the sequence $\{B_i^{(u)}\}$ is i.i.d., and $\mathbb{E}[-B_i^{(u)} \mid \theta^* \in [L_k, U_k], E_{k,p,q}] \leq 0$. By Lemma 1, the probability above is at most $\frac{\delta}{4 \log \frac{1}{2\epsilon}}$.

Thus, $\Pr(\theta^* \notin [L_{k+1}, R_{k+1}] | E_{k,p,q}) \leq \frac{\delta}{2 \log \frac{1}{2\epsilon}}$ when $\theta^* \in [L_k, U_k]$. Similarly, when $\theta^* \in [V_k, R_k]$, we can show $\Pr(\theta^* \notin [L_{k+1}, R_{k+1}] | E_{k,p,q}) \leq \Pr(O_{k,p,q}^{(2)} | E_{k,p,q}) + \Pr(O_{k,p,q}^{(4)} | E_{k,p,q}) \leq \frac{\delta}{2 \log \frac{1}{2\epsilon}}$.

Therefore, $\Pr(\theta^* \notin [L_{k+1}, R_{k+1}] \mid E_{k,p,q}) \leq \frac{\delta}{2 \log \frac{1}{2\epsilon}}$, and thus $\Pr(|\hat{\theta} - \theta^*| > \epsilon) \leq \delta/2$. \square

Proof of Theorem 2. Define T_k to be the number of iterations of the loop at Line 6, $T = \sum_{k=0}^{\log \frac{1}{2\epsilon} - 1} T_k$. For any numbers $m_1, m_2, \dots, m_{\log \frac{1}{2\epsilon} - 1}$, we have:

$$\begin{aligned}
\Pr(T \geq m) &\leq \Pr(|\hat{\theta} - \theta^*| > \epsilon) + \Pr\left(|\hat{\theta} - \theta^*| < \epsilon \text{ and } T \geq \sum_{k=0}^{\log \frac{1}{2\epsilon} - 1} m_k\right) \\
&\leq \frac{\delta}{2} + \Pr\left(T \geq \sum_{k=0}^{\log \frac{1}{2\epsilon} - 1} m_k \text{ and } |\hat{\theta} - \theta^*| < \epsilon\right) \\
&\leq \frac{\delta}{2} + \sum_{k=0}^{\log \frac{1}{2\epsilon} - 1} \Pr(T_k \geq m_k \text{ and } |\hat{\theta} - \theta^*| < \epsilon) \\
&\leq \frac{\delta}{2} + \sum_{k=0}^{\log \frac{1}{2\epsilon} - 1} \Pr(T_k \geq m_k \text{ and } \theta^* \in [L_k, R_k])
\end{aligned} \tag{1}$$

The first and the third inequality follows by union bounds. The second follows by Theorem 1. The last follows since $|\hat{\theta} - \theta^*| < \epsilon$ is equivalent to $\theta^* \in [L_{\log \frac{1}{2\epsilon}}, R_{\log \frac{1}{2\epsilon}}]$, which implies $\theta^* \in [L_k, R_k]$ for all $k = 0, \dots, \log \frac{1}{2\epsilon} - 1$.

We define \mathbb{Q}_k as in the previous proof. For all $k = 0, \dots, \log \frac{1}{2\epsilon} - 1$,

$$\begin{aligned}
&\Pr(T_k \geq m_k \text{ and } \theta^* \in [L_k, R_k]) \\
&= \sum_{(p,q) \in \mathbb{Q}_k: p \leq \theta^* \leq q} \Pr(T_k \geq m_k, L_k = p, R_k = q) \\
&= \sum_{(p,q) \in \mathbb{Q}_k: p \leq \theta^* \leq q} \Pr(T_k \geq m_k \mid L_k = p, R_k = q) \Pr(L_k = p, R_k = q)
\end{aligned}$$

Thus, in order to prove the query complexity of Algorithm 1 is $O\left(\sum_{k=0}^{\log \frac{1}{2\epsilon} - 1} m_k\right)$, it suffices to show that $\Pr(T_k \geq m_k \mid L_k = p, R_k = q) \leq \frac{\delta}{2 \log \frac{1}{2\epsilon}}$ for any $k = 0, \dots, \log \frac{1}{2\epsilon} - 1$, $(p, q) \in \mathbb{Q}_k$ and $p \leq \theta^* \leq q$.

For each k, p, q , define event $E_{k,p,q}$ to be the event $L_k = p, R_k = q$. Define $l_k = q - p = \left(\frac{3}{4}\right)^k$, N_k to be $\tilde{\Theta}\left(\frac{1}{f(l_k/4)} l_k^{-2\beta}\right)$. The logarithm factor of N_k is to be specified later. Define $S_n^{(u)}$ and $S_n^{(v)}$ to be the size of array $B^{(u)}$ and $B^{(v)}$ before Line 16 respectively.

To show $\Pr(T_k \geq N_k \mid E_{k,p,q}) \leq \frac{\delta}{2 \log \frac{1}{2\epsilon}}$, it suffices to show that on event $E_{k,p,q}$, with probability at least $1 - \frac{\delta}{2 \log \frac{1}{2\epsilon}}$, if $n = N_k$ then at least one of the two calls to CheckSignificant between Line 22 and Line 27 will return true.

On event $E_{k,p,q}$, if $\theta^* \in [L_k, M_k]$ (note that on event $E_{k,p,q}$, L_k and M_k are deterministic), then $|V_k - \theta^*| \geq \frac{l_k}{4}$. We will show

$$p_1 := \Pr\left(\text{CheckSignificant}\left(\left\{B_i^{(v)}\right\}_{i=1}^{S_{N_k}^{(v)}}, \frac{\delta}{4 \log \frac{1}{2\epsilon}}\right) \text{ returns false} \mid E_{k,p,q}\right) \leq \frac{\delta}{2 \log \frac{1}{2\epsilon}}$$

To prove this, we will first show that $S_{N_k}^{(v)}$, the length of the array $B^{(v)}$, is large with high probability, and then apply Lemma 2 to show that CheckSignificant will return true if $S_{N_k}^{(v)}$ is large.

By definition, $S_{N_k}^{(v)} = \sum_{i=1}^{N_k} A_i^{(v)}$. By Condition 2, $\mathbb{E}[A_i^{(v)} | E_{k,p,q}] = \Pr(Y \neq \perp | X = V_k, E_{k,p,q}) \geq f\left(\frac{l_k}{4}\right)$.

On event $E_{k,p,q}$, $\{A_i^{(v)}\}$ is a sequence of i.i.d. random variables. By the multiplicative Chernoff bound, $\Pr\left(S_{N_k}^{(v)} \leq \frac{1}{2}N_k f\left(\frac{l_k}{4}\right) \mid E_{k,p,q}\right) \leq \exp\left(-N_k f\left(\frac{l_k}{4}\right)/8\right)$.

Now,

$$p_1 \leq \Pr\left(\text{CheckSignificant}\left(\left\{B_i^{(v)}\right\}_{i=1}^{S_{N_k}^{(v)}}, \frac{\delta}{4 \log \frac{1}{2\epsilon}}\right) \text{ returns false}, S_{N_k}^{(v)} \geq \frac{1}{2}N_k f\left(\frac{l_k}{4}\right) \mid E_{k,p,q}\right) \\ + \Pr\left(S_{N_k}^{(v)} < \frac{1}{2}N_k f\left(\frac{l_k}{4}\right) \mid E_{k,p,q}\right)$$

By Condition 2 and $|V_k - \theta^*| \geq \frac{l_k}{4}$, $\mathbb{E}[B_i^{(v)} | E_{k,p,q}] \geq C\left(\frac{l_k}{4}\right)^\beta$. On event $E_{k,p,q}$, $\{B_i^{(v)}\}$ is a sequence of i.i.d. random variables. Thus, On event $E_{k,p,q}$, by Lemma 2, with probability at least $1 - \frac{\delta}{4 \log \frac{1}{2\epsilon}}$, CheckSignificant will return true if $\frac{1}{2}N_k f\left(\frac{l_k}{4}\right) = \Theta\left(\frac{1}{l_k^{2\beta}} \ln \frac{\ln 1/\epsilon}{\delta} [\ln \ln]_+ \frac{1}{l_k^{2\beta}}\right)$. We have already proved $\Pr\left(S_{N_k}^{(v)} \leq \frac{1}{2}N_k f\left(\frac{l_k}{4}\right) \mid E_{k,p,q}\right) \leq \exp\left(-N_k f\left(\frac{l_k}{4}\right)/8\right)$. By setting $N_k = \Theta\left(\frac{1}{f(l_k/4)} l_k^{-2\beta} \ln \frac{\ln 1/\epsilon}{\delta} [\ln \ln]_+ \frac{1}{l_k^{2\beta}}\right)$, we can ensure p_1 is at most $\delta/2 \log \frac{1}{2\epsilon}$.

Now we have proved on event $E_{k,p,q}$, if $\theta^* \in [L_k, M_k]$, then

$$\Pr\left(\text{CheckSignificant}\left(\left\{B_i^{(v)}\right\}_{i=1}^{S_{N_k}^{(v)}}, \frac{\delta}{4 \log \frac{1}{2\epsilon}}\right) \text{ returns true} \mid E_{k,p,q}\right) \geq 1 - \frac{\delta}{2 \log \frac{1}{2\epsilon}}$$

Likewise, on event $E_{k,p,q}$, if $\theta^* \in [M_k, R_k]$, then

$$\Pr\left(\text{CheckSignificant}\left(\left\{-B_i^{(u)}\right\}_{i=1}^{S_{N_k}^{(u)}}, \frac{\delta}{4 \log \frac{1}{2\epsilon}}\right) \text{ returns true} \mid E_{k,p,q}\right) \geq 1 - \frac{\delta}{2 \log \frac{1}{2\epsilon}}$$

Therefore, we have shown $\Pr(T_k \geq N_k \mid E_{k,p,q}) \leq \frac{\delta}{2 \log \frac{1}{2\epsilon}}$ for any k, p, q . By (1), with probability at least $1 - \delta$, the number of samples queried is at most

$$\sum_{k=0}^{\log \frac{1}{2\epsilon} - 1} O\left(\frac{1}{f\left(\left(\frac{3}{4}\right)^k / 4\right)} \left(\frac{3}{4}\right)^{-2\beta k} \ln \frac{\ln 1/\epsilon}{\delta} [\ln \ln]_+ \left(\frac{3}{4}\right)^{-2\beta k}\right) \\ = O\left(\frac{\epsilon^{-2\beta}}{f(\epsilon/2)} \ln \frac{1}{\epsilon} \left(\ln \frac{1}{\delta} + \ln \ln \frac{1}{\epsilon}\right) [\ln \ln]_+ \frac{1}{\epsilon}\right)$$

□

Proof of Theorem 3. For each k in Algorithm 1 at Line 3, Let $l_k = R_k - L_k$. Let $N_k = \eta \frac{1}{f(l_k/4)} \ln \frac{4 \log \frac{1}{2\epsilon}}{\delta}$, where η is a constant to be specified later. As with the previous proof, it suffices to show $\Pr(T_k \geq N_k \mid E_{k,p,q}) \leq \frac{\delta}{2 \log \frac{1}{2\epsilon}}$ where event $E_{k,p,q}$ is defined to be $L_k = p, R_k = q, T_k$ is the number of iterations at the loop at Line 6.

On event $E_{k,p,q}$, we will show that the loop at Line 6 will terminate after $n = N_k$ with probability at least $1 - \frac{\delta}{2 \log \frac{1}{2\epsilon}}$.

Suppose for now $\theta^* \in [M_k, R_k]$. Let $Z_i = A_i^{(u)} - A_i^{(m)}$, $\zeta = \theta^* - M_k$. Clearly, $|Z_i| \leq 1$. On event $E_{k,p,q}$, sequence $\{Z_i\}$ is i.i.d.. By Condition 3, $\mathbb{E}[Z_i \mid E_{k,p,q}] = f(\zeta + \frac{l_k}{4}) - f(\zeta) \geq$

$cf(\zeta + \frac{l_k}{4})$ since $\zeta \leq \frac{2}{3}(\zeta + \frac{l_k}{4})$. $\text{Var}[Z_i | E_{k,p,q}] = \text{Var}[A_i^{(u)} | E_{k,p,q}] + \text{Var}[A_i^{(m)} | E_{k,p,q}] \stackrel{(a)}{\leq} \mathbb{E}[A_i^{(u)} | E_{k,p,q}] + \mathbb{E}[A_i^{(m)} | E_{k,p,q}] = f(\zeta + \frac{l_k}{4}) + f(\zeta) \stackrel{(b)}{\leq} 2f(\zeta + \frac{l_k}{4})$ where (a) follows by $A_i \in \{0, 1\}$ and (b) follows by the monotonicity of f . Thus, on event $E_{k,p,q}$, by Lemma 4, if we set η sufficiently large (independent of l_k, ϵ, δ), then with probability at least $1 - \frac{\delta}{4 \log \frac{1}{2\epsilon}}$ $\text{CheckSignificant-Var}(\{Z_i\}_{i=1}^{N_k}, \frac{\delta}{4 \log \frac{1}{2\epsilon}})$ in Procedure 2 returns true.

Similarly, we can show that on event $E_{k,p,q}$, if $\theta^* \in [L_k, M_k]$, by Lemma 4, with probability at least $1 - \frac{\delta}{4 \log \frac{1}{2\epsilon}}$, $\text{CheckSignificant-Var}(\{A_i^{(v)} - A_i^{(m)}\}_{i=1}^{N_k}, \frac{\delta}{4 \log \frac{1}{2\epsilon}})$ returns true.

Therefore, the loop at Line 6 will terminate after $n = N_k$ with probability at least $1 - \frac{\delta}{4 \log \frac{1}{2\epsilon}}$ on event $E_{k,p,q}$. Therefore, with probability at least $1 - \delta$, the number of samples queried is at most $\sum_{k=0}^{\log \frac{1}{2\epsilon} - 1} \frac{1}{f((\frac{3}{4})^k / 4)} \ln \frac{\ln 1/\epsilon}{\delta} = O\left(\frac{1}{f(\epsilon/2)} \ln \frac{1}{\epsilon} \left(\ln \frac{1}{\delta} + \ln \ln \frac{1}{\epsilon}\right)\right)$. \square

A.3 The d-dimensional case

To prove the d -dimensional case, we only need to use a union bound to show that with high probability all calls of Algorithm 1 succeed, and consequently the output boundary g produced by polynomial interpolation is close to the true underlying boundary due to the smoothness assumption of g^* .

Proof of Theorem 8. For $q \in \left\{0, 1, \dots, \frac{M}{\gamma} - 1\right\}^{d-1}$, define the “polynomial interpolation” version of g^* as

$$g_q^*(\tilde{\mathbf{x}}) = \sum_{l \in I_q \cap \mathcal{L}} g^*(l) Q_{q,l}(\tilde{\mathbf{x}})$$

Recall that we choose $M = O(\epsilon^{-1/\gamma})$.

By Theorem 1, each run of Algorithm 1 at the line 3 of Algorithm 3 will return a g_l such that $|g_l - g_q^*(l)| \leq \epsilon$ with probability at least $1 - \delta/2M^{d-1}$.

$$\begin{aligned} & \|g - g^*\| \\ &= \sum_{q \in \{0, \dots, M/\gamma - 1\}^{d-1}} \|(g_q - g^*) \mathbb{1}\{\tilde{\mathbf{x}} \in I_q\}\| \\ &\leq \sum_{q \in \{0, \dots, M/\gamma - 1\}^{d-1}} \|(g_q - g_q^*) \mathbb{1}\{\tilde{\mathbf{x}} \in I_q\}\| + \|(g_q^* - g^*) \mathbb{1}\{\tilde{\mathbf{x}} \in I_q\}\| \end{aligned}$$

$$\begin{aligned} \|(g_q^* - g^*) \mathbb{1}\{\tilde{\mathbf{x}} \in I_q\}\| &= \int_{I_q} |g_q^*(\tilde{\mathbf{x}}) - g^*(\tilde{\mathbf{x}})| d\tilde{\mathbf{x}} \\ &= O\left(\int_{I_q} M^{-\gamma} d\tilde{\mathbf{x}}\right) \\ &= O(M^{-\gamma-d+1}) \end{aligned}$$

The second equality follows from Lemma 3 of [6] that $|g_q(\tilde{\mathbf{x}}) - g^*(\tilde{\mathbf{x}})| = O(M^{-\gamma})$ since g^* is γ -Hölder smooth.

$$\begin{aligned}
& \| (g_q - g_q^*) \mathbb{1}\{\tilde{\mathbf{x}} \in I_q\} \| \\
&= \sum_{l \in I_q \cap \mathcal{L}} |g_l - g_q^*(l)| \|Q_{q,l}\| \\
&\leq \sum_{l \in I_q \cap \mathcal{L}} \epsilon \|Q_q\| \\
&= O(\epsilon M^{-d+1})
\end{aligned}$$

Therefore, overall we have $\|g - g^*\| \leq O(M^{-\gamma-d+1} + \epsilon M^{-d+1}) \left(\frac{M}{\gamma}\right)^{d-1} = O(\epsilon)$. \square

Proof of Theorem 9. By Theorem 2, each run of Algorithm 1 at the line 3 of Algorithm 3 will make $\tilde{O}\left(\frac{d}{f(\epsilon/2)}\epsilon^{-2\beta}\right)$ queries with probability at least $1 - \delta/M^{d-1}$, thus by a union bound, the total number of queries made is $\tilde{O}\left(\frac{d}{f(\epsilon/2)}\epsilon^{-2\beta-\frac{d-1}{\gamma}}\right)$ with probability at least $1 - \delta$. \square

Proof of Theorem 10. The proof is similar to the previous proof. \square

B Proof of lower bounds

First, we introduce some notations for this section. Given a labeler L and an active learning algorithm \mathcal{A} , denote by $P_{L,\mathcal{A}}^n$ the distribution of n samples $\{(X_i, Y_i)\}_{i=1}^n$ where Y_i is drawn from distribution $P_L(Y|X_i)$ and X_i is drawn by the active learning algorithm based solely on the knowledge of $\{(X_j, Y_j)\}_{j=1}^{i-1}$. We will drop the subscripts from $P_{L,\mathcal{A}}^n$ and $P_{\mathcal{L}}(Y|X)$ when it is clear from the context. For a sequence $\{X_i\}_{i=1}^\infty$ denote by X^n the subsequence $\{X_1, \dots, X_n\}$.

Definition 1. For any distributions P, Q on a countable support, define KL-divergence as $d_{\text{KL}}(P, Q) = \sum_x P(x) \ln \frac{P(x)}{Q(x)}$. For two random variables X, Y , define the mutual information as $I(X; Y) = d_{\text{KL}}(P(X, Y) \| P(X)P(Y))$.

We will use Fano's method shown as below to prove the lower bounds.

Lemma 5. Let Θ be a class of parameters, and $\{P_\theta : \theta \in \Theta\}$ be a class of probability distributions indexed by Θ over some sample space \mathcal{X} . Let $d : \Theta \times \Theta \rightarrow \mathbb{R}$ be a semi-metric. Let $\mathcal{V} = \{\theta_1, \dots, \theta_M\} \subseteq \Theta$ such that $\forall i \neq j, d(\theta_i, \theta_j) \geq 2s > 0$. Let $\bar{P} = \frac{1}{M} \sum_{\theta \in \mathcal{V}} P_\theta$. If $d_{\text{KL}}(P_\theta \| \bar{P}) \leq \delta$ for any $\theta \in \mathcal{V}$, then for any algorithm $\hat{\theta}$ that given a sample X drawn from P_θ outputs $\hat{\theta}(X) \in \Theta$, the following inequality holds:

$$\sup_{\theta \in \Theta} P_\theta \left(d(\theta, \hat{\theta}(X)) \geq s \right) \geq 1 - \frac{\delta + \ln 2}{\ln M}$$

Proof. For any algorithm $\hat{\theta}$, define a test function $\hat{\Psi} : \mathcal{X} \rightarrow \{1, \dots, M\}$ such that $\hat{\Psi}(X) = \arg \min_{i \in \{1, \dots, M\}} d(\hat{\theta}(X), \theta_i)$. We have

$$\sup_{\theta \in \Theta} P_\theta \left(d(\theta, \hat{\theta}(X)) \geq s \right) \geq \max_{\theta \in \mathcal{V}} P_\theta \left(d(\theta, \hat{\theta}(X)) \geq s \right) \geq \max_{i \in \{1, \dots, M\}} P_{\theta_i} \left(\hat{\Psi}(X) \neq i \right)$$

Let V be a random variable uniformly taking values from \mathcal{V} , and X be drawn from P_V . By Fano's Inequality, for any test function $\Psi : \mathcal{X} \rightarrow \{1, \dots, M\}$

$$\max_{i \in \{1, \dots, M\}} P_{\theta_i} (\Psi(X) \neq i) \geq 1 - \frac{I(V; X) + \ln 2}{\ln M}$$

The desired result follows by the fact that $I(V; X) = \frac{1}{M} \sum_{\theta \in \mathcal{V}} d_{\text{KL}}(P_\theta \| \bar{P})$. \square

B.1 The one dimensional case

*Proof of Theorem 5.*² Without lose of generality, let $C = C' = 1$ (C is defined in Condition 2). Let $\epsilon \leq \frac{1}{4} \min \left\{ \left(\frac{1}{2} \right)^{1/\beta}, \left(\frac{4}{5} \right)^{1/\alpha}, \frac{1}{4} \right\}$. We will prove the desired result using Lemma 5.

First, we construct \mathcal{V} and P_θ . For any $k \in \{0, 1, 2, 3\}$, let $P_{L_k}(Y | X)$ be the distribution of the labeler L_k 's response with the ground truth $\theta_k = k\epsilon$:

$$\begin{aligned} P_{L_k}(Y = \perp | x) &= 1 - \left| x - \frac{1}{2} - k\epsilon \right|^\alpha \\ P_{L_k}(Y = 0 | x) &= \begin{cases} (x - \frac{1}{2} - k\epsilon)^\alpha \left(1 - (x - \frac{1}{2} - k\epsilon)^\beta \right) / 2 & x > \frac{1}{2} + k\epsilon \\ (\frac{1}{2} + k\epsilon - x)^\alpha \left(1 + (\frac{1}{2} + k\epsilon - x)^\beta \right) / 2 & x \leq \frac{1}{2} + k\epsilon \end{cases} \\ P_{L_k}(Y = 1 | x) &= \begin{cases} (x - \frac{1}{2} - k\epsilon)^\alpha \left(1 + (x - \frac{1}{2} - k\epsilon)^\beta \right) / 2 & x > \frac{1}{2} + k\epsilon \\ (\frac{1}{2} + k\epsilon - x)^\alpha \left(1 - (\frac{1}{2} + k\epsilon - x)^\beta \right) / 2 & x \leq \frac{1}{2} + k\epsilon \end{cases} \end{aligned}$$

Clearly, P_{L_k} complies with Conditions 1, 2 and 3.

Define P_k^n to be the distribution of n samples $\{(X_i, Y_i)\}_{i=1}^n$ where Y_i is drawn from distribution $P_{L_k}(Y | X_i)$ and X_i is drawn by the active learning algorithm based solely on the knowledge of $\{(X_j, Y_j)\}_{j=1}^{i-1}$.

Define $\bar{P}_L = \frac{1}{4} \sum_j P_{L_j}$ and $\bar{P}^n = \frac{1}{4} \sum_j P_k^n$. We take Θ to be $[0, 1]$, and $d(\theta_1, \theta_2) = |\theta_1 - \theta_2|$ in Lemma 5. To use Lemma 5, we need to bound $d_{\text{KL}}(P_k^n \| \bar{P}^n)$ for $k \in \{0, 1, 2, 3\}$.

For any $k \in \{0, 1, 2, 3\}$,

$$\begin{aligned} & d_{\text{KL}}(P_k^n \| \bar{P}^n) \\ &= \mathbb{E}_{P_k^n} \left(\ln \frac{P_k^n(\{(X_i, Y_i)\}_{i=1}^n)}{\bar{P}^n(\{(X_i, Y_i)\}_{i=1}^n)} \right) \\ &= \mathbb{E}_{P_k^n} \left(\ln \frac{P_k^n(X_1) P_k^n(Y_1 | X_1) P_k^n(X_2 | X_1, Y_1) \cdots P_k^n(Y_n | X_1, Y_1, \dots, X_n)}{\bar{P}^n(X_1) \bar{P}^n(Y_1 | X_1) \bar{P}^n(X_2 | X_1, Y_1) \cdots \bar{P}^n(Y_n | X_1, Y_1, \dots, X_n)} \right) \\ &\stackrel{(a)}{=} \mathbb{E}_{P_k^n} \left(\ln \frac{\prod_{i=1}^n P_{L_k}(Y_i | X_i)}{\prod_{i=1}^n \bar{P}_L(Y_i | X_i)} \right) \\ &= \sum_{i=1}^n \mathbb{E}_{P_k^n} \left(\mathbb{E}_{P_k^n} \left(\ln \frac{P_{L_k}(Y_i | X_i)}{\bar{P}_L(Y_i | X_i)} \mid X^n \right) \right) \\ &\leq n \max_{x \in [0, 1]} d_{\text{KL}}(P_{L_k}(Y | x) \| \bar{P}_L(Y | x)) \end{aligned} \tag{2}$$

(a) follows by the fact that $P_k^n(X_{i+1} | X_1, Y_1, \dots, X_i, Y_i) = \bar{P}^n(X_{i+1} | X_1, Y_1, \dots, X_i, Y_i)$ since X_{i+1} is drawn by the same active learning algorithm based solely on the knowledge of $\{(X_j, Y_j)\}_{j=1}^i$ regardless of the labeler's response distribution, and the fact that $P_k^n(Y_i | X_1, Y_1, \dots, X_i) = P_{L_k}(Y_i | X_i)$ and $\bar{P}^n(Y_i | X_1, Y_1, \dots, X_i) = \bar{P}_L(Y_i | X_i)$ by definition.

For any $k \in \{1, 2, 3\}$, $x \in [0, 1]$,

$$\bar{P}_L(\cdot | x) \geq \frac{P_{L_0}(\cdot | x) + P_{L_k}(\cdot | x)}{4} \tag{3}$$

For any $k \in \{0, 1, 2, 3\}$, $x \in [0, 1]$, $y \in \{1, -1, \perp\}$

²Actually we can use Le Cam's method to prove this one dimensional case (which only needs to construct 2 distributions instead of 4 here), but this proof can be generalized to the multidimensional case more easily.

$$\begin{aligned}
& (\bar{P}_L(Y = y | x) - P_{L_k}(Y = y | x))^2 \\
&= \left(\sum_j \frac{1}{4} (P_{L_j}(Y = y | x) - P_{L_0}(Y = y | x)) + (P_{L_0}(Y = y | x) - P_{L_k}(Y = y | x)) \right)^2 \\
&\leq \left(\frac{5}{16} \sum_{j>0} (P_{L_j}(Y = y | x) - P_{L_0}(Y = y | x))^2 + 5 (P_{L_0}(Y = y | x) - P_{L_k}(Y = y | x))^2 \right) \\
&\leq 6 \sum_{j>0} (P_{L_j}(Y = y | x) - P_{L_0}(Y = y | x))^2 \tag{4}
\end{aligned}$$

where the first inequality follows by $\left(\sum_{i=0}^4 a_i\right)^2 \leq 5 \sum_{i=0}^4 a_i^2$ by letting $a_j = \frac{1}{4} (P_{L_j}(Y = y | x) - P_{L_0}(Y = y | x))$ for $j = 0, \dots, 3$ and $a_4 = P_{L_0}(Y = y | x) - P_{L_k}(Y = y | x)$, and noting that $a_0 = 0$ under this setting.

Thus,

$$\begin{aligned}
& d_{\text{KL}}(P_{L_k}(Y | x) \| \bar{P}_L(Y | x)) \\
&\leq \sum_y \frac{1}{\bar{P}_L(Y = y | x)} (P_{L_k}(Y = y | x) - \bar{P}_L(Y = y | x))^2 \\
&\leq 24 \sum_{j>0} \sum_y \frac{1}{P_{L_j}(y | x) + P_{L_0}(y | x)} (P_{L_j}(Y = y | x) - P_{L_0}(Y = y | x))^2 \\
&\leq O(\epsilon^\alpha)
\end{aligned}$$

The first inequality follows from Lemma 10. The second inequality follows by (3) and (4). The last inequality follows by applying Lemma 11 to $P_{L_0}(\cdot | x)$ and $P_{L_j}(\cdot | x)$ and the assumption $\alpha \leq 2$.

Therefore, we have $d_{\text{KL}}(P_k^n \| \bar{P}_0^n) = nO(\epsilon^\alpha)$. By setting $n = \epsilon^{-\alpha}$, we get $d_{\text{KL}}(P_k^n \| \bar{P}_0^n) \leq O(1)$, and thus by Lemma 5,

$$\sup_{\theta} P_{\theta} \left(d(\theta, \hat{\theta}(X)) \geq \Omega(\epsilon) \right) \geq 1 - \frac{O(1) + \ln 2}{\ln 4} = O(1)$$

□

B.2 The d-dimensional case

Again, we will use Lemma 5 to prove the lower bounds for d -dimensional cases. We first construct $\{P_{\theta} : \theta \in \Theta\}$ using a similar idea with [6], and then use Lemma 12 to select a subset $\tilde{\Theta} \subset \Theta$ to apply Lemma 5.

Proof of Theorem 6. Again, without lose of generality, let $C = 1$. Recall that for $\mathbf{x} = (x_1, \dots, x_d) \in \mathbb{R}^d$, we have defined $\tilde{\mathbf{x}}$ to be (x_1, \dots, x_{d-1}) . Define $m = (\frac{1}{\epsilon})^{1/\gamma}$. $\mathcal{L} = \{0, \frac{1}{m}, \dots, \frac{m-1}{m}\}^{d-1}$, $h(\tilde{\mathbf{x}}) = \prod_{i=1}^{d-1} \exp\left(-\frac{1}{1-4x_i^2}\right) \mathbb{1}\{|x_i| < \frac{1}{2}\}$, $\phi_l(\tilde{\mathbf{x}}) = Km^{-\gamma}h(m(\tilde{\mathbf{x}}-l)-\frac{1}{2})$ where $l \in \mathcal{L}$. It is easy to check $\phi_l(\tilde{\mathbf{x}})$ is (K, γ) -Hölder smooth and has bounded support $[l_1, l_1 + \frac{1}{m}] \times \dots \times [l_{d-1}, l_{d-1} + \frac{1}{m}]$, which implies that for different $l_1, l_2 \in \mathcal{L}$, the support of ϕ_{l_1} and ϕ_{l_2} do not intersect.

Let $\Omega = \{0, 1\}^{m^{d-1}}$. For any $\omega \in \Omega$, define $g_{\omega}(\tilde{\mathbf{x}}) = \sum_{l \in \mathcal{L}} \omega_l \phi_l(\tilde{\mathbf{x}})$. For each $\omega \in \Omega$, define the conditional distribution of labeler L_{ω} 's response as follows:

For $x_d \leq A$, $P_{L_{\omega}}(y = \perp | \mathbf{x}) = 1 - f(A)$, $P_{L_{\omega}}(y \neq \perp | \mathbf{x}, y \neq \perp) = \frac{1}{2} \left(1 - |x_d - g_{\omega}(\tilde{\mathbf{x}})|^{\beta}\right)$;

For $x_d \geq A$, $P_{L_\omega}(y = \perp | \mathbf{x}) = 1 - f(x_d)$, $P_{L_\omega}(y \neq \mathbb{I}(x_d > g_\omega(\tilde{\mathbf{x}})) | \mathbf{x}, y \neq \perp) = \frac{1}{2} (1 - x_d^\beta)$.

Here, $A = c \max \phi(\tilde{\mathbf{x}}) = c' \epsilon$ for some constants c, c' .

It can be easily verified that P_{L_ω} satisfies Conditions 1 and 2. Note that $g_\omega(\tilde{\mathbf{x}})$ can be seen as the underlying decision boundary for labeler P_{L_ω} .

Define P_ω^n to be the distribution of n samples $\{(X_i, Y_i)\}_{i=1}^n$ where Y_i is drawn from distribution $P_{L_\omega}(Y | X_i)$ and X_i is drawn by the active learning algorithm based solely on the knowledge of $\{(X_j, Y_j)\}_{j=1}^{i-1}$.

By Lemma 12, when ϵ is small enough so that m^{d-1} is large enough, there is a subset $\{\omega^{(1)}, \dots, \omega^{(M)}\} \subset \Omega$ such that $\|\omega^{(i)} - \omega^{(j)}\|_0 \geq m^{d-1}/12$ for any $0 \leq i < j \leq M$ and $M \geq 2^{m^{d-1}/48}$. Define $P_i^n = P_{\omega^{(i)}}^n$, $\bar{P}^n = \frac{1}{M} \sum_{i=1}^M P_i^n$.

Next, we will apply Lemma 5 to $\{\omega^{(1)}, \dots, \omega^{(M)}\}$ with $d(\omega^{(i)}, \omega^{(j)}) = \|g_{\omega^{(i)}} - g_{\omega^{(j)}}\|$. We will lower-bound $d(\omega^{(i)}, \omega^{(j)})$ and upper-bound $d_{\text{KL}}(P_i^n \| \bar{P}^n)$.

For any $1 \leq i < j \leq M$,

$$\begin{aligned} & \|g_{\omega^{(i)}} - g_{\omega^{(j)}}\| \\ &= \sum_{l \in \{1, \dots, m\}^{d-1}} \left| \omega_l^{(i)} - \omega_l^{(j)} \right| K m^{-\gamma-(d-1)} \|h\| \\ &\geq m^{d-1}/12 * K m^{-\gamma-(d-1)} \|h\| \\ &= K m^{-\gamma} \|h\| / 12 \\ &= \Theta(\epsilon) \end{aligned}$$

By the convexity of KL-divergence, $d_{\text{KL}}(P_i^n \| \bar{P}^n) \leq \frac{1}{M} \sum_{j=1}^M d_{\text{KL}}(P_i^n \| P_j^n)$, so it suffices to upper-bound $d_{\text{KL}}(P_i^n \| P_j^n)$ for any i, j .

For any $1 < i, j \leq M$,

$$\begin{aligned} & d_{\text{KL}}(P_i^n \| P_j^n) \\ &\leq n \max_{\mathbf{x} \in [0,1]^d} d_{\text{KL}}(P_{L_{\omega^{(i)}}}^n(Y | \mathbf{x}) \| P_{L_{\omega^{(j)}}}^n(Y | \mathbf{x})) \\ &= n \max_{\mathbf{x} \in [0,1]^d} P_{L_{\omega^{(i)}}}^n(Y \neq \perp | \mathbf{x}) d_{\text{KL}}(P_{L_{\omega^{(i)}}}^n(Y | \mathbf{x}, Y \neq \perp) \| P_{L_{\omega^{(j)}}}^n(Y | \mathbf{x}, Y \neq \perp)) \end{aligned}$$

The inequality follows as (2) in the proof of Theorem 5. The equality follows since $P_\omega(y = \perp | \mathbf{x})$ is the same for all $\omega \in \Omega$.

If $x_d \geq A$, then $P_{L_{\omega^{(i)}}}^n(Y | \mathbf{x}, Y \neq \perp) = P_{L_{\omega^{(j)}}}^n(Y | \mathbf{x}, Y \neq \perp)$, so $d_{\text{KL}}(P_{L_{\omega^{(i)}}}^n(Y | \mathbf{x}, Y \neq \perp) \| P_{L_{\omega^{(j)}}}^n(Y | \mathbf{x}, Y \neq \perp)) = 0$. If $x_d < A$, then $P_{L_{\omega^{(i)}}}^n(Y \neq \perp | \mathbf{x}) = f(A)$. Therefore,

$$d_{\text{KL}}(P_i^n \| P_j^n) \leq n f(A) \max_{\mathbf{x} \in [0,1]^d} d_{\text{KL}}(P_{L_{\omega^{(i)}}}^n(Y | \mathbf{x}, Y \neq \perp) \| P_{L_{\omega^{(j)}}}^n(Y | \mathbf{x}, Y \neq \perp))$$

Apply Lemma 10 to $P_{L_{\omega^{(i)}}}^n(Y | \mathbf{x}, Y \neq \perp)$ and $P_{L_{\omega^{(j)}}}^n(Y | \mathbf{x}, Y \neq \perp)$, and noting they are bounded above by a constant, we have $\max_{\mathbf{x} \in [0,1]^d} d_{\text{KL}}(P_{L_{\omega^{(i)}}}^n(Y | \mathbf{x}, Y \neq \perp) \| P_{L_{\omega^{(j)}}}^n(Y | \mathbf{x}, Y \neq \perp)) = O(A^{2\beta})$. Thus,

$$d_{\text{KL}}(P_i^n \| P_j^n) \leq n f(A) O(A^{2\beta}) = n f(c' \epsilon) O(\epsilon^{2\beta})$$

By setting $n = \frac{1}{f(c' \epsilon)} \epsilon^{-2\beta - \frac{d-1}{\gamma}}$, we get $d_{\text{KL}}(P_i^n \| P_j^n) \leq O(\epsilon^{-\frac{d-1}{\gamma}})$. The desired results follows by Lemma 5. \square

The proof of Theorem 7 follows the same structure.

Proof of Theorem 7. As in the proof of Theorem 6, let $C = C' = 1$, and define $m = (\frac{1}{\epsilon})^{1/\gamma}$. $\mathcal{L} = \{0, \frac{1}{m}, \dots, \frac{m-1}{m}\}^{d-1}$, $h(\tilde{\mathbf{x}}) = \prod_{i=1}^{d-1} \exp\left(-\frac{1}{1-4x_i^2}\right) \mathbb{1}\{|x_i| < \frac{1}{2}\}$, $\phi_l(\tilde{\mathbf{x}}) = Km^{-\gamma}h(m(\tilde{\mathbf{x}} - l) - \frac{1}{2})$ where $l \in \mathcal{L}$. Let $\Omega = \{0, 1\}^{m^{d-1}}$. For any $\omega \in \Omega$, define $g_\omega(\tilde{\mathbf{x}}) = \frac{1}{2} + \sum_{l \in \mathcal{L}} \omega_l \phi_l(\tilde{\mathbf{x}})$, which can be seen as a decision boundary. $A = \max \phi(\tilde{\mathbf{x}}) = c'\epsilon$ for some constants c' .

Let $g_+(\tilde{\mathbf{x}}) = g_{(1,1,\dots,1)}(\tilde{\mathbf{x}}) = \sum_{l \in \mathcal{L}} \phi_l(\tilde{\mathbf{x}})$, $g_-(\tilde{\mathbf{x}}) = g_{(0,0,\dots,0)}(\tilde{\mathbf{x}}) = 0$. In other words, g_+ is the “highest” boundary, and g_- is the “lowest” boundary.

For each $\omega \in \Omega$, define the conditional distribution of labeler L_ω ’s response as follows:

$$P_{L_\omega}(y = \pm | \mathbf{x}) = 1 - |x_d - g_\omega(\tilde{\mathbf{x}})|^\alpha$$

$$P_{L_\omega}(y \neq \mathbb{I}(x_d > g_\omega(\tilde{\mathbf{x}})) | \mathbf{x}, y \neq \pm) = \frac{1}{2} \left(1 - |x_d - g_\omega(\tilde{\mathbf{x}})|^\beta\right)$$

It can be easily verified that P_{L_ω} satisfies Conditions 1, 2, and 3.

Let $P_+(\cdot | \mathbf{x}) = P_{L_{(1,1,\dots,1)}}(\cdot | \mathbf{x})$, $P_-(\cdot | \mathbf{x}) = P_{L_{(0,0,\dots,0)}}(\cdot | \mathbf{x})$. By the construction of g , for any $\mathbf{x} \in [0, 1]^d$, any $\omega \in \Omega$, $P_{L_\omega}(\cdot | \mathbf{x})$ equals either $P_+(\cdot | \mathbf{x})$ or $P_-(\cdot | \mathbf{x})$.

Define P_ω^n to be the distribution of n samples $\{(X_i, Y_i)\}_{i=1}^n$ where Y_i is drawn from distribution $P_{L_\omega}(Y | X_i)$ and X_i is drawn by the active learning algorithm based solely on the knowledge of $\{(X_j, Y_j)\}_{j=1}^{i-1}$.

By Lemma 12, when ϵ is small enough so that m^{d-1} is large enough, there is a subset $\Omega' = \{\omega^{(1)}, \dots, \omega^{(M)}\} \subset \Omega$ such that (i) (well-separated) $\|\omega^{(i)} - \omega^{(j)}\|_0 \geq m^{d-1}/12$ for any $0 \leq i < j \leq M$, $M \geq 2^{m^{d-1}/48}$; and (ii) (well-balanced) for any $j = 1, \dots, m^{d-1}$, $\frac{1}{24} \leq \frac{1}{M} \sum_{i=1}^M \omega_j^{(i)} \leq \frac{3}{24}$.

Define $P_i^n = P_{\omega^{(i)}}^n$, $\bar{P}^n = \frac{1}{M} \sum_{i=1}^M P_i^n$. Define $P_{L_i} = P_{L_{\omega^{(i)}}}$, $\bar{P}_L = \frac{1}{M} \sum_{i=1}^M P_{L_i}$. By the well-balanced property, for any $\mathbf{x} \in [0, 1]^d$, $\bar{P}_L(\cdot | \mathbf{x})$ is between $\frac{1}{24}P_+(\cdot | \mathbf{x}) + \frac{23}{24}P_-(\cdot | \mathbf{x})$ and $\frac{3}{24}P_+(\cdot | \mathbf{x}) + \frac{21}{24}P_-(\cdot | \mathbf{x})$. Therefore

$$\bar{P}_L(\cdot | \mathbf{x}) \geq \frac{1}{24} (P_+(\cdot | \mathbf{x}) + P_-(\cdot | \mathbf{x})) \quad (5)$$

Moreover, since $P_{L_i}(\cdot | \mathbf{x})$ can only take $P_+(\cdot | \mathbf{x})$ or $P_-(\cdot | \mathbf{x})$ for any \mathbf{x} ,

$$|P_{L_i}(\cdot | \mathbf{x}) - \bar{P}_L(\cdot | \mathbf{x})| \leq |P_+(\cdot | \mathbf{x}) - P_-(\cdot | \mathbf{x})| \quad (6)$$

Next, we will apply Lemma 5 to $\{\omega^{(1)}, \dots, \omega^{(M)}\}$ with $d(\omega^{(i)}, \omega^{(j)}) = \|g_{\omega^{(i)}} - g_{\omega^{(j)}}\|$. We already know from the proof of Theorem 6 $\|g_{\omega^{(i)}} - g_{\omega^{(j)}}\| = \Omega(\epsilon)$.

For any $0 < i \leq M$, $d_{\text{KL}}(P_i^n \| \bar{P}_0^n) \leq n \max_{\mathbf{x} \in [0, 1]^d} d_{\text{KL}}(P_{L_i}(Y | \mathbf{x}) \| \bar{P}_L(Y | \mathbf{x}))$. For any $\mathbf{x} \in [0, 1]^d$,

$$\begin{aligned} & d_{\text{KL}}(P_{L_i}(Y | \mathbf{x}) \| \bar{P}_L(Y | \mathbf{x})) \\ & \leq \sum_y \frac{1}{\bar{P}_L(Y = y | \mathbf{x})} (P_{L_i}(Y = y | \mathbf{x}) - \bar{P}_L(Y = y | \mathbf{x}))^2 \\ & \leq \sum_y \frac{24}{P_+(y | \mathbf{x}) + P_-(y | \mathbf{x})} (P_+(Y = y | \mathbf{x}) - P_-(Y = y | \mathbf{x}))^2 \\ & \leq O(A^\alpha) \end{aligned}$$

The first inequality follows from Lemma 10. The second inequality follows by (5) and (6). The last inequality follows by applying Lemma 11 to $P_+(\cdot | \mathbf{x})$ and $P_-(\cdot | \mathbf{x})$, setting the ϵ in Lemma 11 to be $g_\omega(\tilde{\mathbf{x}})$, and using $g_\omega(\tilde{\mathbf{x}}) \leq A$ and the assumption $\alpha \leq 2$.

Therefore, we have

$$d_{\text{KL}}(P_i^n \parallel P_0^n) \leq nO(A^\alpha) = nO(\epsilon^\alpha)$$

By setting $n = \epsilon^{-\alpha - \frac{d-1}{\gamma}}$, we get $d_{\text{KL}}(P_i^n \parallel P_0^n) \leq O\left(\epsilon^{-\frac{d-1}{\gamma}}\right)$. Thus by Lemma 5,

$$\sup_{\theta} P_{\theta} \left(d(\theta, \hat{\theta}(X)) \geq \Omega(\epsilon) \right) \geq 1 - \frac{O\left(\epsilon^{-\frac{d-1}{\gamma}}\right) + \ln 2}{\epsilon^{-\frac{d-1}{\gamma}}/48} = O(1)$$

, from which the desired result follows. \square

C Technical lemmas

C.1 Concentration bounds

In this subsection, we define Y_1, Y_2, \dots to be a sequence of i.i.d. random variables. Assume $Y_1 \in [-2, 2]$, $\mathbb{E}Y_1 = 0$, $\text{Var}(Y_1) = \sigma^2 \leq 4$. Define $V_n = \frac{n}{n-1} \left(\sum_{i=1}^n Y_i^2 - \frac{1}{n} \left(\sum_{i=1}^n Y_i \right)^2 \right)$. It is easy to check $\mathbb{E}V_n = n\sigma^2$.

We need following two results from [21]

Lemma 6. ([21], Theorem 2) *Take any $0 < \delta < 1$. Then there is an absolute constant D_0 such that with probability at least $1 - \delta$, for all n simultaneously,*

$$\left| \sum_{i=1}^n Y_i \right| \leq D_0 \left(1 + \ln \frac{1}{\delta} + \sqrt{n\sigma^2 [\ln \ln]_+ (n\sigma^2) + n\sigma^2 \ln \frac{1}{\delta}} \right)$$

Lemma 7. ([21], Lemma 3) *Take any $0 < \delta < 1$. Then there is an absolute constant K_0 such that with probability at least $1 - \delta$, for all n simultaneously,*

$$n\sigma^2 \leq K_0 \left(1 + \ln \frac{1}{\delta} + \sum_{i=1}^n Y_i^2 \right)$$

We note that Proposition 1 is immediate from Lemma 6 since $\text{Var}(Y_i) \leq 4$.

Lemma 8. *Take any $0 < \delta < 1$. Then there is an absolute constant K_3 such that with probability at least $1 - \delta$, for all $n \geq \ln \frac{1}{\delta}$ simultaneously,*

$$n\sigma^2 \leq K_3 \left(1 + \ln \frac{1}{\delta} + V_n \right)$$

Proof. By Lemma 7, with probability at least $1 - \delta/2$, for all n ,

$$n\sigma^2 \leq K_0 \left(\sum_{i=1}^n Y_i^2 + \ln \frac{2}{\delta} + 1 \right) = K_0 \left(\frac{n-1}{n} V_n + \frac{1}{n} \left(\sum_{i=1}^n Y_i \right)^2 + \ln \frac{2}{\delta} + 1 \right)$$

By Lemma 6, with probability at least $1 - \delta/2$, for all n ,

$$\begin{aligned}
\frac{1}{n} \left(\sum_{i=1}^n Y_i \right)^2 &< \frac{1}{n} \left(D_0 \left(1 + \ln \frac{2}{\delta} + \sqrt{n\sigma^2 [\ln \ln]_+ (n\sigma^2) + n\sigma^2 \ln \frac{2}{\delta}} \right) \right)^2 \\
&= \frac{D_0^2}{n} \left(1 + \ln \frac{2}{\delta} \right)^2 + D_0^2 \sigma^2 [\ln \ln]_+ (n\sigma^2) + D_0^2 \sigma^2 \ln \frac{2}{\delta} \\
&\quad + 2D_0^2 \left(1 + \ln \frac{2}{\delta} \right) \sqrt{\frac{\sigma^2 [\ln \ln]_+ (n\sigma^2) + \sigma^2 \ln \frac{2}{\delta}}{n}} \\
&\leq K_1 \left(1 + \ln \frac{1}{\delta} + [\ln \ln]_+ (n\sigma^2) \right)
\end{aligned}$$

for some absolute constant K_1 . The last inequality follows by $n \geq \ln \frac{1}{\delta}$.

Thus, by a union bound, with probability at least $1 - \delta$, for all n , $n\sigma^2 \leq K_0 V_n + K_0(K_1 + 2) \ln \frac{1}{\delta} + K_0 K_1 [\ln \ln]_+ (n\sigma^2) + K_0(K_1 + 3)$.

Let $K_2 > 0$ be an absolute constant such that $\forall x \geq K_2$, $K_0 K_1 [\ln \ln]_+ x \leq \frac{x}{2}$.

Now if $n\sigma^2 \geq K_2$, then $n\sigma^2 \leq K_0 V_n + K_0(K_1 + 2) \ln \frac{1}{\delta} + \frac{n\sigma^2}{2} + K_0(K_1 + 3)$, and thus

$$n\sigma^2 \leq 2K_0 V_n + 2K_0(K_1 + 2) \ln \frac{1}{\delta} + 2K_0(K_1 + 3) + K_2 \quad (7)$$

If $n\sigma^2 \leq K_2$, clearly (7) holds. This concludes the proof. \square

We note that Proposition 2 is immediate by applying above lemma to Lemma 6.

Lemma 9. *Take any $\delta, n > 0$. Then with probability at least $1 - \delta$,*

$$V_n \leq 4n\sigma^2 + 8 \ln \frac{1}{\delta}$$

Proof. Applying Bernstein's Inequality to Y_i^2 , and noting that $\text{Var}(Y_i^2) \leq 4\sigma^2$ since $|Y_i| \leq 2$, we have with probability at least $1 - \delta$,

$$\begin{aligned}
\sum_{i=1}^n Y_i^2 &\leq \frac{4}{3} \ln \frac{1}{\delta} + n\sigma^2 + \sqrt{8n\sigma^2 \ln \frac{1}{\delta}} \\
&\leq 4 \ln \frac{1}{\delta} + 2n\sigma^2
\end{aligned}$$

The last inequality follows by the fact that $\sqrt{4ab} \leq a + b$.

The desired result follows by noting that $V_n = \frac{n}{n-1} \left(\sum_{i=1}^n Y_i^2 - \frac{1}{n} \left(\sum_{i=1}^n Y_i \right)^2 \right) \leq 2 \sum_{i=1}^n Y_i^2$. \square

C.2 Bounds of distances among probability distributions

Lemma 10. *If P, Q are two probability distributions on a countable support \mathcal{X} , then*

$$d_{\text{KL}}(P \parallel Q) \leq \sum_x \frac{(P(x) - Q(x))^2}{Q(x)}$$

Proof.

$$\begin{aligned}
d_{\text{KL}}(P \parallel Q) &= \sum_x P(x) \ln \frac{P(x)}{Q(x)} \\
&\leq \sum_x P(x) \left(\frac{P(x)}{Q(x)} - 1 \right) \\
&= \sum_x \frac{(P(x) - Q(x))^2}{Q(x)}
\end{aligned}$$

The first inequality follows by $\ln x \leq x - 1$. The second equality follows by $\sum_x P(x) \left(\frac{P(x)}{Q(x)} - 1 \right) = \sum_x \left(\frac{P^2(x) - P(x)Q(x)}{Q(x)} - P(x) + Q(x) \right) = \sum_x \frac{(P(x) - Q(x))^2}{Q(x)}$. \square

Define

$$\begin{aligned}
P_0(Y = \perp | x) &= 1 - \left| x - \frac{1}{2} \right|^\alpha \\
P_0(Y = 0 | x) &= \begin{cases} \left(x - \frac{1}{2} \right)^\alpha \left(1 - \left(x - \frac{1}{2} \right)^\beta \right) / 2 & x > \frac{1}{2} \\ \left(\frac{1}{2} - x \right)^\alpha \left(1 + \left(\frac{1}{2} - x \right)^\beta \right) / 2 & x \leq \frac{1}{2} \end{cases} \\
P_0(Y = 1 | x) &= \begin{cases} \left(x - \frac{1}{2} \right)^\alpha \left(1 + \left(x - \frac{1}{2} \right)^\beta \right) / 2 & x > \frac{1}{2} \\ \left(\frac{1}{2} - x \right)^\alpha \left(1 - \left(\frac{1}{2} - x \right)^\beta \right) / 2 & x \leq \frac{1}{2} \end{cases}
\end{aligned}$$

and

$$\begin{aligned}
P_1(Y = \perp | x) &= 1 - \left| x - \epsilon - \frac{1}{2} \right|^\alpha \\
P_1(Y = 0 | x) &= \begin{cases} \left(x - \epsilon - \frac{1}{2} \right)^\alpha \left(1 - \left(x - \epsilon - \frac{1}{2} \right)^\beta \right) / 2 & x > \epsilon + \frac{1}{2} \\ \left(\epsilon + \frac{1}{2} - x \right)^\alpha \left(1 + \left(\epsilon + \frac{1}{2} - x \right)^\beta \right) / 2 & x \leq \epsilon + \frac{1}{2} \end{cases} \\
P_1(Y = 1 | x) &= \begin{cases} \left(x - \epsilon - \frac{1}{2} \right)^\alpha \left(1 + \left(x - \epsilon - \frac{1}{2} \right)^\beta \right) / 2 & x > \epsilon + \frac{1}{2} \\ \left(\epsilon + \frac{1}{2} - x \right)^\alpha \left(1 - \left(\epsilon + \frac{1}{2} - x \right)^\beta \right) / 2 & x \leq \epsilon + \frac{1}{2} \end{cases}
\end{aligned}$$

Lemma 11. *Let P_0, P_1 be the distributions defined above. If $x \in [0, 1]$, $\epsilon \leq \min \left\{ \left(\frac{1}{2} \right)^{1/\beta}, \left(\frac{4}{5} \right)^{1/\alpha}, \frac{1}{4} \right\}$, then*

$$\sum_y \frac{(P_0(Y = y|x) - P_1(Y = y|x))^2}{P_0(Y = y|x) + P_1(Y = y|x)} = O(\epsilon^\alpha + \epsilon^2) \quad (8)$$

Proof. By symmetry, it suffices to show for $0 \leq x \leq \frac{1+\epsilon}{2}$. Let $t = \frac{1}{2} + \epsilon - x$.

We first show (8) holds for $\frac{\epsilon}{2} \leq t \leq \epsilon$ (i.e. $\frac{1}{2} \leq x \leq \frac{1+\epsilon}{2}$).

We claim $\min_y (P_0(Y = y|X = t) + P_1(Y = y|X = t)) \geq \frac{1}{2} \left(\frac{\epsilon}{2} \right)^\alpha$. This is because:

- $P_0(Y = \perp | X = t) + P_1(Y = \perp | X = t) = 1 - (\epsilon - t)^\alpha + 1 - t^\alpha \geq 2 - 2\epsilon^\alpha \geq \frac{1}{2} \left(\frac{\epsilon}{2} \right)^\alpha$
where the last inequality follows by $\epsilon \leq \left(\frac{4}{5} \right)^{1/\alpha}$;
- $2(P_0(Y = 0|X = t) + P_1(Y = 0|X = t)) = (\epsilon - t)^\alpha \left(1 - (\epsilon - t)^\beta \right) + t^\alpha (1 + t^\beta) \geq t^\alpha (1 + t^\beta) \geq \left(\frac{\epsilon}{2} \right)^\alpha$. Therefore, $P_0(Y = 0|X = t) + P_1(Y = 0|X = t) \geq \frac{1}{2} \left(\frac{\epsilon}{2} \right)^\alpha$.

- Similarly, $P_0(Y = 1|X = t) + P_1(Y = 1|X = t) \geq \frac{1}{2} \left(\frac{\epsilon}{2}\right)^\alpha$.

Besides,

$$\begin{aligned}
& \sum_y (P_0(Y = y|X = t) - P_1(Y = y|X = t))^2 \\
&= (t^\alpha - (\epsilon - t)^\alpha)^2 + \frac{1}{4} \left(t^\alpha (1 - t^\beta) - (\epsilon - t)^\alpha (1 + (\epsilon - t)^\beta) \right)^2 \\
&\quad + \frac{1}{4} \left(t^\alpha (1 + t^\beta) - (\epsilon - t)^\alpha (1 - (\epsilon - t)^\beta) \right)^2 \\
&= (t^\alpha - (\epsilon - t)^\alpha)^2 + \frac{1}{4} \left(t^\alpha - (\epsilon - t)^\alpha - t^{\alpha+\beta} - (\epsilon - t)^{\alpha+\beta} \right)^2 \\
&\quad + \frac{1}{4} \left(t^\alpha - (\epsilon - t)^\alpha + t^{\alpha+\beta} + (\epsilon - t)^{\alpha+\beta} \right)^2 \\
&\stackrel{(a)}{\leq} (t^\alpha - (\epsilon - t)^\alpha)^2 + \frac{1}{2} (t^\alpha - (\epsilon - t)^\alpha)^2 + \frac{1}{2} \left(t^{\alpha+\beta} + (\epsilon - t)^{\alpha+\beta} \right)^2 \\
&\quad + \frac{1}{2} (t^\alpha - (\epsilon - t)^\alpha)^2 + \frac{1}{2} \left(t^{\alpha+\beta} + (\epsilon - t)^{\alpha+\beta} \right)^2 \\
&= 2(t^\alpha - (\epsilon - t)^\alpha)^2 + \left(t^{\alpha+\beta} + (\epsilon - t)^{\alpha+\beta} \right)^2 \\
&\leq 2\epsilon^{2\alpha} + 4\epsilon^{2\alpha+2\beta} \\
&\leq 6\epsilon^{2\alpha}
\end{aligned}$$

where (a) follows by the inequality $(a + b)^2 \leq 2a^2 + 2b^2$ for any a, b .

Therefore, we get $\sum_y \frac{(P_0(Y=y|x) - P_1(Y=y|x))^2}{P_0(Y=y|x) + P_1(Y=y|x)} \leq \frac{\sum_y (P_0(Y=y|x) - P_1(Y=y|x))^2}{\min_y (P_0(Y=y|x) + P_1(Y=y|x))} \leq 12 * 2^\alpha \epsilon^\alpha$ when $\frac{1}{2} \leq x \leq \frac{1+\epsilon}{2}$.

Next, We show (8) holds for $\epsilon \leq t \leq \frac{1}{2} + \epsilon$ (i.e. $0 \leq x \leq \frac{1}{2}$). We will show $\frac{(P_0(Y=y|x) - P_1(Y=y|x))^2}{P_0(Y=y|x) + P_1(Y=y|x)} = O(\epsilon^\alpha + \epsilon^2)$ for $Y = \perp, 1, 0$.

For $Y = \perp$, for the denominator,

$$P_0(Y = \perp | X = t) + P_1(Y = \perp | X = t) = 2 - t^\alpha - (t - \epsilon)^\alpha \geq 2 - \left(\frac{3}{4}\right)^\alpha - \left(\frac{1}{2}\right)^\alpha$$

For the numerator,

$$(P_0(Y = \perp | X = t) - P_1(Y = \perp | X = t))^2 = (t^\alpha - (t - \epsilon)^\alpha)^2 = t^{2\alpha} \left(1 - \left(1 - \frac{\epsilon}{t}\right)^\alpha\right)^2$$

By Lemma 13, if $\alpha \geq 1$, $t^{2\alpha} \left(1 - \left(1 - \frac{\epsilon}{t}\right)^\alpha\right)^2 \leq t^{2\alpha} \left(\alpha \frac{\epsilon}{t}\right)^2 = t^{2\alpha-2} (\alpha\epsilon)^2 = O(\epsilon^2)$. If $0 \leq \alpha \leq 1$, $t^{2\alpha} \left(1 - \left(1 - \frac{\epsilon}{t}\right)^\alpha\right)^2 \leq t^{2\alpha} \left(\frac{\epsilon}{t}\right)^2 = t^{2\alpha-2} \epsilon^2 \leq \epsilon^{2\alpha}$.

Thus, we have $\frac{(P_0(Y=\perp|x) - P_1(Y=\perp|x))^2}{P_0(Y=\perp|x) + P_1(Y=\perp|x)} = O(\epsilon^{2\alpha} + \epsilon^2)$.

For $Y = 1$, for the denominator,

$$\begin{aligned}
2(P_0(Y = 1|X = t) + P_1(Y = 1|X = t)) &= t^\alpha (1 - t^\beta) + (t - \epsilon)^\alpha (1 - (t - \epsilon)^\beta) \\
&\geq t^\alpha (1 - t^\beta) \\
&\geq t^\alpha \left(1 - \left(\frac{3}{4}\right)^\beta\right)
\end{aligned}$$

For the numerator,

$$\begin{aligned}
& (P_0(Y=1|X=t) - P_1(Y=1|X=t))^2 \\
&= \frac{1}{4} \left(t^\alpha (1-t^\beta) - (t-\epsilon)^\alpha (1-(t-\epsilon)^\beta) \right)^2 \\
&\leq \frac{1}{2} (t^\alpha - (t-\epsilon)^\alpha)^2 + \frac{1}{2} (t^{\alpha+\beta} - (t-\epsilon)^{\alpha+\beta})^2 \\
&= \frac{1}{2} t^{2\alpha} \left(1 - (1 - \frac{\epsilon}{t})^\alpha \right)^2 + \frac{1}{2} t^{2\alpha+2\beta} \left(1 - (1 - \frac{\epsilon}{t})^{\alpha+\beta} \right)^2 \\
&\leq \frac{1}{2} t^{2\alpha} \left(1 - (1 - \frac{\epsilon}{t})^\alpha \right)^2 + \frac{1}{2} t^{2\alpha} \left(1 - (1 - \frac{\epsilon}{t})^{\alpha+\beta} \right)^2
\end{aligned}$$

If $\alpha \geq 1$, by Lemma 13, $\frac{1}{2} t^{2\alpha} (1 - (1 - \frac{\epsilon}{t})^\alpha)^2 + \frac{1}{2} t^{2\alpha} (1 - (1 - \frac{\epsilon}{t})^{\alpha+\beta})^2 \leq \frac{1}{2} t^{2\alpha} (\alpha \frac{\epsilon}{t})^2 + \frac{1}{2} t^{2\alpha} ((\alpha + \beta) \frac{\epsilon}{t})^2 = \left(\frac{1}{2} \alpha^2 + \frac{1}{2} (\alpha + \beta)^2 \right) t^{2\alpha-2} \epsilon^2$. Thus, $\frac{(P_0(Y=1|x) - P_1(Y=1|x))^2}{P_0(Y=1|x) + P_1(Y=1|x)} \leq \left(\frac{1}{2} \alpha^2 + \frac{1}{2} (\alpha + \beta)^2 \right) t^{\alpha-2} \epsilon^2 / \left(1 - (\frac{3}{4})^\beta \right)$ which is $O(\epsilon^2)$ if $\alpha \geq 2$ and $O(\epsilon^\alpha)$ if $\alpha \leq 2$.

If $\alpha \leq 1$ and $\alpha + \beta \geq 1$, by Lemma 13, $\frac{1}{2} t^{2\alpha} (1 - (1 - \frac{\epsilon}{t})^\alpha)^2 + \frac{1}{2} t^{2\alpha} (1 - (1 - \frac{\epsilon}{t})^{\alpha+\beta})^2 \leq \frac{1}{2} t^{2\alpha} (\frac{\epsilon}{t})^2 + \frac{1}{2} t^{2\alpha} ((\alpha + \beta) \frac{\epsilon}{t})^2 = \left(\frac{1}{2} + \frac{1}{2} (\alpha + \beta)^2 \right) t^{2\alpha-2} \epsilon^2 \leq \left(\frac{1}{2} + \frac{1}{2} (\alpha + \beta)^2 \right) t^{2\alpha-2} \epsilon^2$. Thus, $\frac{(P_0(Y=1|x) - P_1(Y=1|x))^2}{P_0(Y=1|x) + P_1(Y=1|x)} \leq \left(\frac{1}{2} + \frac{1}{2} (\alpha + \beta)^2 \right) t^{\alpha-2} \epsilon^2 / \left(1 - (\frac{3}{4})^\beta \right) = O(\epsilon^\alpha)$.

If $\alpha \leq 1, \alpha + \beta \leq 1$, by Lemma 13, $\frac{1}{2} t^{2\alpha} (1 - (1 - \frac{\epsilon}{t})^\alpha)^2 + \frac{1}{2} t^{2\alpha} (1 - (1 - \frac{\epsilon}{t})^{\alpha+\beta})^2 \leq \frac{1}{2} t^{2\alpha} (\frac{\epsilon}{t})^2 + \frac{1}{2} t^{2\alpha} (\frac{\epsilon}{t})^2 = t^{2\alpha-2} \epsilon^2$. Thus, $\frac{(P_0(Y=1|x) - P_1(Y=1|x))^2}{P_0(Y=1|x) + P_1(Y=1|x)} \leq t^{\alpha-2} \epsilon^2 / \left(1 - (\frac{3}{4})^\beta \right) = O(\epsilon^\alpha)$.

Therefore, we have $\frac{(P_0(Y=1|x) - P_1(Y=1|x))^2}{P_0(Y=1|x) + P_1(Y=1|x)} = O(\epsilon^\alpha + \epsilon^2)$.

Likewise, we can get $\frac{(P_0(Y=0|x) - P_1(Y=0|x))^2}{P_0(Y=0|x) + P_1(Y=0|x)} = O(\epsilon^\alpha + \epsilon^2)$. So we prove $\sum_y \frac{(P_0(Y=y|x) - P_1(Y=y|x))^2}{P_0(Y=y|x) + P_1(Y=y|x)} = O(\epsilon^\alpha + \epsilon^2)$ when $x \leq \frac{1}{2}$. This concludes the proof. \square

C.3 Other lemmas

Lemma 12. ([20], Lemma 4) For sufficiently large $d > 0$, there is a subset $M \subset \{0, 1\}^d$ with following properties: (i) $|M| \geq 2^{d/48}$; (ii) $\|v - v'\|_0 > \frac{d}{12}$ for any two distinct $v, v' \in M$; (iii) for any $i = 1, \dots, d$, $\frac{1}{24} \leq \frac{1}{|M|} \sum_{v \in M} v_i \leq \frac{3}{24}$.

Lemma 13. If $x \leq 1, r \geq 1$, then $(1-x)^r \geq 1-rx$ and $1 - (1-x)^r \leq rx$.

If $0 \leq x \leq 1, 0 \leq r \leq 1$, then $(1-x)^r \geq \frac{1-x}{1-rx}$ and $1 - (1-x)^r \leq \frac{rx}{1-(1-r)x} \leq x$.

Inequalities above are known as Bernoulli's inequalities. One proof can be found in [16].

Lemma 14. Suppose ϵ, τ are positive numbers and $\delta \leq \frac{1}{2}$. Suppose $\{Z_i\}_{i=1}^\infty$ is a sequence of i.i.d random variables bounded by 1, $\mathbb{E}Z_i \geq \tau\epsilon$, and $\text{Var}(Z_i) = \sigma^2 \leq 2\epsilon$. Define $V_n = \frac{n}{n-1} \left(\sum_{i=1}^n Z_i - \frac{1}{n} \left(\sum_{i=1}^n Z_i \right)^2 \right)$, $q_n = q(n, V_n, \delta)$ as Procedure 2. If $n \geq \frac{n}{\tau\epsilon} \ln \frac{1}{\delta}$ for some sufficiently large number η (to be specified in the proof), then with probability at least $1 - \delta$, $\frac{q_n}{n} - \mathbb{E}Z_i \leq -\tau\epsilon/2$.

Proof. By Lemma 9, with probability at least $1 - \delta$, $V_n \leq 4n\sigma^2 + 8 \ln \frac{1}{\delta}$, which implies

$$q_n \leq D_1 \left(1 + \ln \frac{1}{\delta} + \sqrt{\left(4n\sigma^2 + 9 \ln \frac{1}{\delta} + 1 \right) \left([\ln \ln]_+ (4n\sigma^2 + 9 \ln \frac{1}{\delta} + 1) + \ln \frac{1}{\delta} \right)} \right)$$

We denote the RHS by q .

On this event, we have

$$\begin{aligned}
\frac{q_n}{n} - \mathbb{E}Z_i &\leq \frac{q}{n} - \tau\epsilon \\
&= \tau\epsilon \left(\frac{q}{n\tau\epsilon} - 1 \right) \\
&\stackrel{(a)}{\leq} \tau\epsilon \left(\frac{2D_1}{\eta} + \frac{D_1}{\eta \ln \frac{1}{\delta}} \sqrt{\frac{9\eta}{\tau} \ln \frac{1}{\delta} \left([\ln \ln]_+ \left(\frac{9\eta}{\tau} \ln \frac{1}{\delta} \right) + \ln \frac{1}{\delta} \right) - 1} \right) \\
&= \tau\epsilon \left(\frac{2D_1}{\eta} + D_1 \sqrt{\frac{9}{\eta\tau \ln \frac{1}{\delta}} [\ln \ln]_+ \left(\frac{9\eta}{\tau} \ln \frac{1}{\delta} \right) + \frac{9}{\eta\tau} - 1} \right)
\end{aligned}$$

where (a) follows from $\frac{q}{n}$ being monotonically decreasing with respect to n . By choosing η sufficiently large, we have $\frac{2D_1}{\eta} + D_1 \sqrt{\frac{9}{\eta\tau \ln \frac{1}{\delta}} [\ln \ln]_+ \left(\frac{9\eta}{\tau} \ln \frac{1}{\delta} \right) + \frac{9}{\eta\tau} - 1} \leq -\frac{1}{2}$, and thus $\frac{q_n}{n} - \mathbb{E}Z_i \leq -\tau\epsilon/2$. \square