Supplementary Materials for Complexity of Inference in Latent Dirichlet Allocation

A Proof of Lemma 2

Proof of Lemma 2. Assume there are T sets each having $k \ge 3$ elements, and let Φ be the optimal LDA objective. Define $F(n) = \log \Gamma(n + \alpha)$. Since l_{it} is constant across all topics, the linear term in Eq. 2 will be a constant K. First, note that, if there is a perfect matching,

$$\Phi \ge \frac{n}{k} F(k) + (T - \frac{n}{k}) F(0) + K.$$
(16)

The F(0) term is the contribution of unused topics. Otherwise, assume that the best packing has $\gamma \leq cn/k$ sets, each with k elements. Then, by the properties of the log-gamma function,

$$\Phi \le \gamma F(k) + \frac{n - \gamma k}{k - 1} F(k - 1) + (T - \frac{n}{k}) F(0) + K,$$
(17)

where we assume, conservatively, that all of the remaining words are explained by topics assigned (k-1) words. Also, since there was no perfect matching, there were at most $T - \frac{n}{k}$ unused topics. Using our bound on γ , we have

$$\Phi \le \frac{cn}{k}F(k) + \frac{n - \frac{cn}{k}k}{k - 1}F(k - 1) + (T - \frac{n}{k})F(0) + K$$
(18)

$$= \frac{cn}{k}F(k) + \frac{n(1-c)}{k-1}F(k-1) + (T-\frac{n}{k})F(0) + K$$
(19)

$$= \frac{dn}{k}F(k) + (T - \frac{n}{k})F(0) + K,$$
 (20)

where

$$d := c + (1 - c)\beta, \quad \text{for } \beta := \frac{k}{F(k)} \frac{F(k - 1)}{k - 1}.$$
 (21)

Note that $F(k)/k \to \infty$ as $k \to \infty$. Along with the convexity of F, it follows that there exists a k_0 such that $\beta < 1$ for all $k > k_0$. Note that $k > (3 + \alpha)^2$ suffices. This implies that d < 1, which shows that there is a non-zero gap between the possible values of Φ . \Box

Note that the maximum concentration objective, $F(n) = n \log n$, satisfies the conditions on F and, in particular, we have $\beta < 1$ for k = 3.

B Derivation of MAP θ objective

$$\Pr(\theta|\mathbf{w}) \propto \sum_{z_1,\dots,z_N} \Pr(\theta) \Pr(w_1,\dots,w_N,z_1,\dots,z_N|\theta)$$
(22)

$$= \Pr(\theta) \sum_{\substack{z_1, \dots, z_N \\ N}} \prod_{i=1}^N \Pr(z_i | \theta) \Pr(w_i | z_i)$$
(23)

$$= \Pr(\theta) \prod_{i=1}^{N} \sum_{z_i} \Pr(z_i|\theta) \Pr(w_i|z_i)$$
(24)

$$= \operatorname{Pr}(\theta) \prod_{i=1}^{N} \sum_{t=1}^{T} \theta_t \operatorname{Pr}(w_i | z_i = t)$$
(25)

$$\propto \prod_{t=1}^{T} \theta_t^{\alpha_t - 1} \prod_{i=1}^{N} \sum_{t=1}^{T} \theta_t \Pr(w_i | z_i = t).$$
(26)

C Proof of Lemma 3

If $\epsilon \geq K(\alpha, T, N)$ the claim trivially holds. Assume for the purpose of contradiction that there exists a word \hat{i} such that $\theta_{\hat{t}}^* < K(\alpha, T, N)$, where $\hat{t} = \arg \max_t \psi_{\hat{i}t} \theta_t^*$.

Let Y denote the set of topics $t \neq \hat{t}$ such that $\theta_t^* \geq 2\epsilon$. Let $\beta_1 = \sum_{t \in Y} \theta_t^*$ and $\beta_2 = \sum_{t \notin Y, t \neq \hat{t}} \theta_t^*$. Note that $\beta_2 < 2T\epsilon$. Consider $\hat{\theta}$ defined as follows:

$$\hat{\theta}_{\hat{t}} = \frac{1}{N} \tag{27}$$

$$\hat{\theta}_t = \left(\frac{1-\beta_2 - \frac{1}{N}}{\beta_1}\right) \theta_t^* \text{ for } t \in Y$$
(28)

$$\hat{\theta}_t = \theta_t^* \text{ for } t \notin Y, t \neq \hat{t}.$$
 (29)

Note that this construction implies the bound $\hat{\theta}_t \ge (1 - 2T\epsilon - \frac{1}{N}) \theta_t^*$ for $t \in Y$. Assuming $n \ge 4$ and $\epsilon \le \frac{1}{2TN}$, we have that $\hat{\theta}_t \ge \frac{1}{2}\theta_t^* \ge \epsilon$ for $t \in Y$, so $\hat{\theta}$ is feasible.

We will show that $\Phi(\hat{\theta}) > \Phi(\theta^*)$, contradicting the optimality of θ^* . First we need the following upper bound, which uses the fact that $\theta_{\hat{t}}^* < \frac{1}{N}$:

$$\frac{1 - \beta_2 - \frac{1}{n}}{\beta_1} = \frac{1 - \beta_2 - \frac{1}{N}}{1 - \beta_2 - \theta_i^*}$$
(30)

Then, we have:

$$\frac{P(\hat{\theta})}{\alpha - 1} = \sum_{t=1}^{T} \log(\hat{\theta}_t)$$
(32)

$$= \log \frac{1}{N} + \sum_{t \in Y} \log \left(\frac{1 - \beta_2 - \frac{1}{N}}{\beta_1} \right) + \sum_{t \neq \hat{t}} \log \theta_t^*$$
(33)

$$\leq \log \frac{1}{N} + \sum_{t \neq \hat{t}} \log \theta_t^*.$$
(34)

Thus,

$$\frac{P(\hat{\theta}) - P(\theta^*)}{\alpha - 1} \le \log \frac{1}{N} - \log \theta_{\hat{t}}^* \tag{35}$$

which when $\alpha < 1$ gives the inequality:

$$P(\hat{\theta}) - P(\theta^*) \geq (\alpha - 1) \left(\log \frac{1}{N} - \log \theta_{\hat{t}}^* \right)$$
(36)

$$= (1-\alpha) \Big(\log N + \log \theta_{\hat{t}}^* \Big). \tag{37}$$

Moving on to the second term, we have:

$$L(\hat{\theta}) = \sum_{j \in [N]: j \neq \hat{i}} \log\left(\sum_{t} \psi_{jt} \hat{\theta}_{t}\right) + \log\left(\sum_{t} \psi_{\hat{i}t} \hat{\theta}_{t}\right)$$
(38)

$$\geq (N-1)\log\left(\frac{1-\beta_2-\frac{1}{N}}{\beta_1}\right) + \sum_{j\in[N]:\,j\neq\hat{i}}\log\left(\sum_t\psi_{jt}\theta_t^*\right) + \log\left(\frac{\psi_{\hat{i}\hat{t}}}{N}\right) \quad (39)$$

$$\geq (N-1)\log\left(1-\frac{2}{N}\right) + \sum_{j\in[N]:\,j\neq\hat{i}}\log\left(\sum_{t}\psi_{jt}\theta_{t}^{*}\right) + \log\left(\frac{\psi_{\hat{i}\hat{t}}}{N}\right).$$
(40)

$$L(\hat{\theta}) - L(\theta^*) \geq (N-1)\log\left(1 - \frac{2}{N}\right) + \log\left(\frac{\psi_{\hat{i}\hat{t}}}{N}\right) - \log\left(\sum_t \psi_{\hat{i}t}\theta_t^*\right)$$
(41)

$$\geq (N-1)\log\left(1-\frac{2}{N}\right) + \log\left(\frac{\psi_{\hat{i}\hat{t}}}{N}\right) - \log\left(T\psi_{\hat{i}\hat{t}}\theta_{\hat{t}}^*\right) \tag{42}$$

$$= (N-1)\left(\log(N-2) - \log N\right) + \log\left(\frac{1}{N}\right) - \log\left(T\theta_{\hat{t}}^*\right)$$
(43)

$$\geq (N-1)\left(-\frac{2}{N-2}\right) + \log\left(\frac{1}{N}\right) - \log\left(T\theta_{\hat{t}}^*\right) \tag{44}$$

$$\geq -3 + \log\left(\frac{1}{N}\right) - \log\left(T\theta_{\hat{t}}^*\right). \tag{45}$$

where we used the lower bound $\log(N-2) \ge \log N - \frac{2}{N-2}$ that arises from the convexity of the $\log(x)$ function, and again assumed $N \ge 4$.

Finally, putting these two together, we have:

$$\Phi(\hat{\theta}) - \Phi(\theta^*) \ge -3 - \alpha \log N - \log T - \alpha \log \theta_{\hat{t}}^*$$
(46)

Plugging in $\theta_{\hat{t}}^* < K(\alpha, T, N)$ results in $\Phi(\hat{\theta}) - \Phi(\theta^*) > 0$, giving the contradiction.

D Proof of Theorem 8

Here we reduce from the *unique* set cover problem, where we are guaranteed that there is only one minimal size set that covers all elements. It can be shown that Unique Set Cover is NP-hard (under randomized reductions) by using standard reductions from Unique SAT to Vertex Cover, and then from Vertex Cover to Set Cover.

Consider a Unique Set Cover instance and our standard reduction to an LDA instance as described in earlier sections. In particular, let $\mathbf{w} = (w_1, \ldots, w_N)$ denote a Unique Set Cover reduction instance and let $C \subseteq [T]$ denote the unique minimum cover. Let S_i be those sets (topics) that cover element w_i . We will show that for sufficiently small hyperparameters α_t , we can determine whether a set (topic) $t \in C$ by testing the value of $\mathbb{E}[\theta_t|X]$, thus proving that the computation of the latter is NP-hard.

We have

$$p(\theta \mid X) \propto \prod_{t} \theta_t^{\alpha_t - 1} \prod_{i} \sum_{t' \in S_i} \theta_{t'}$$
(47)

$$=\sum_{r}\prod_{t}\theta_{t}^{\alpha_{t}-1+\eta_{t}(r)},\tag{48}$$

where the final summation is over elements $r \in \mathcal{R} := S_1 \times \cdots \times S_N$ and $\eta_t(r) := |\{i : r_i = t\}|$. For $r \in \mathcal{R}$, we write |r| to denote the number of topics t such that $\eta_t(r) \neq 0$.

Let \mathcal{N} denote the set of sequences $n = (n_1, \ldots, n_T)$ such that $n_t \ge 0$ and $\sum_t n_t = N$. For $n \in \mathcal{N}$, define

$$Z(n) := \int \cdots \int \prod_{t} \theta_t^{\alpha_t - 1 + n_t} d\theta_1 \cdots d\theta_T = \frac{\prod_t \Gamma(\alpha_t + n_t)}{\Gamma(\bar{\alpha} + N)},$$
(49)

where $\bar{\alpha} = \sum_t \alpha_t$. It follows that

$$\mathbb{E}[\theta_t|X] = \int \cdots \int \theta_t \cdot p(\theta \mid X) d\theta_1 \cdots d\theta_T$$
(50)

$$= \frac{1}{\sum_{r} Z(\eta(r))} \int \cdots \int \theta_t \sum_{r} \prod_{\tau} \theta_{\tau}^{\alpha_{\tau} - 1 + n_{\tau}(r)} d\theta_1 \cdots d\theta_T$$
(51)

$$=\frac{1}{\sum_{r}Z(\eta(r))}\sum_{r}Z(\eta(r,t)),\tag{52}$$

where $\eta(r,t) \in \mathcal{N}$ is given by $\eta(r,t)_{\tau} = \eta_{\tau}(r) + 1(\tau = t)$. By the identity $\Gamma(z+1) = z\Gamma(z)$, we have $Z(\eta(r,t)) = Z(\eta(r)) \frac{\alpha_t + n_t(r)}{\overline{\alpha} + N}$, and so it follows that

$$\mathbb{E}[\theta_t|X] = \sum_{r \in \mathcal{R}} \bar{Z}(\eta(r)) \frac{\alpha_t + n_t(r)}{\bar{\alpha} + N},$$
(53)

where

$$\bar{Z}(n) := \frac{Z(n)}{\sum_{r' \in \mathcal{R}} Z(\eta(r'))}.$$
(54)

For $c \in [T]$, let \mathcal{R}_c denote the set of those $r \in \mathcal{R}$ such that |r| = c. Recall that $C \subseteq [T]$ is the unique minimum set cover associated with the reduction X. Thus, for $t \in C$,

$$\mathbb{E}[\theta_t|X] = \sum_{r \in \mathcal{R}_{|C|}} \bar{Z}(\eta(r)) \frac{\alpha_t + n_t(r)}{\bar{\alpha} + N} + \sum_{r \in \mathcal{R} \setminus \mathcal{R}_{|C|}} \bar{Z}(\eta(r)) \frac{\alpha_t + n_t(r)}{\bar{\alpha} + N}$$
(55)

$$\geq \frac{\alpha_t + 1}{\bar{\alpha} + N} \sum_{r \in \mathcal{R}_{|C|}} \bar{Z}(\eta(r)), \tag{56}$$

where we have used the observation that $\eta_t(r) \geq 1$ for $r \in \mathcal{R}_{|C|}$. Let $\beta = \sum_{r \in \mathcal{R}_{|C|}} \overline{Z}(\eta(r))$. If $t \notin C$, then $n_t(r) = 0$ for $r \in \mathcal{R}_{|C|}$ and so we have $\mathbb{E}[\theta_t|X] \leq \beta \frac{\alpha_t}{\overline{\alpha}+N} + (1-\beta)$. It follows that if

$$\beta > \frac{1}{2} \left(1 + \frac{\bar{\alpha} + N}{\bar{\alpha} + N + 1} \right) \tag{57}$$

then the minimum cover C contains a topic t if and only if $\mathbb{E}[\theta|X] \geq \frac{1}{4} \left(1 + 3\frac{\bar{\alpha}+N}{\bar{\alpha}+N+1}\right) \frac{\alpha_t+1}{\bar{\alpha}+N}$, and, moreover, we can determine the minimum cover from a bound on β and polynomial approximations to the marginal distributions of the components of θ . We will take $\alpha_t = \alpha$ henceforth, and show that for α small enough, the bound (57) indeed holds.

Let $n, n' \in \mathcal{N}$ be topic counts associated with the minimal cover and some non-minimal cover, respectively. That is, let $n = \eta(r)$ for some $r \in \mathcal{R}_{|C|}$ and let $n' = \eta(r')$ for some $r' \in \mathcal{R}_k$ and k > |C|. We will bound Z(n')/Z(n) in order to bound β . We have

$$\prod_{t} \Gamma(\alpha + n_t) \ge \Gamma(\alpha)^{T - |C|} \Gamma(\alpha + 1)^{|C|},$$
(58)

whereas

$$\prod_{t} \Gamma(\alpha + n'_{t}) \leq \Gamma(\alpha)^{T - |C| - 1} \Gamma(\alpha + 1)^{|C|} \Gamma(\alpha + N - |C|).$$
(59)

Therefore,

$$\frac{Z(n')}{Z(n)} \le \frac{\Gamma(\alpha)^{T-|C|-1} \Gamma(\alpha+1)^{|C|} \Gamma(\alpha+N-|C|)}{\Gamma(\alpha)^{T-|C|} \Gamma(\alpha+1)^{|C|}}$$
(60)

$$=\frac{\Gamma(\alpha+N-|C|)}{\Gamma(\alpha)}.$$
(61)

By the convexity of $\Gamma(1/c)$ in c, we have $\Gamma(\alpha) \ge \alpha^{-1} - \gamma$, where $\gamma \approx .577$ is the Euler constant. Therefore,

$$\frac{Z(n')}{Z(n)} \le \frac{\Gamma(\alpha + N - 1)}{\alpha^{-1} - \gamma} =: \kappa(\alpha).$$
(62)

Then by conservatively assuming that there is only one responsibility corresponding to the minimum cover, we have that

$$\beta \ge \frac{Z(n)}{Z(n) + T^N \kappa(\alpha) Z(n)}.$$
(63)

Therefore, the bound (57) is achieved when

$$\kappa(\alpha) \le \frac{1}{T^N(2\bar{\alpha} + 2N + 1)}.\tag{64}$$

In particular, when $N,T\geq 2$ and

$$\alpha^{-1} > 2T^N \Gamma(N)(2N+2) \tag{65}$$

the marginal expectations can be used to read off the unique minimal set cover.